Data-Driven Program Analysis

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Research Areas

- Program Analysis derives specifications from code
- Program Synthesis derives code from specifications

Program Analysis

- Predict program behavior automatically
 - static or dynamic: before execution at compile-time / at runtime
 - automatic: sw is analyzed by sw ("program analyzers")
- Applications
 - **bug-finding:** e.g., find runtime failures of programs
 - security: e.g., is this app malicious or benign?
 - verification: e.g., does the program meet its specification?
 - compiler optimization: e.g., automatic parallelization

Program Synthesis

- Generate program code from specifications automatically
 - specification: logics, examples, implementation, etc
 - **automatic**: sw is generated by sw ("program synthesizers")
- Applications
 - programming assistance: e.g., complete tricky parts of programs
 - end-user programming: e.g., automate repetitive tasks
 - algorithm discovery: find a new solution for a problem
 - program optimization: find a more efficient implementation
 - automatic patch generation: automatically fix software bugs













Challenge in Static Analysis



Challenge in Static Analysis



Flow-Sensitivity





precise but costly

Flow-Insensitivity



 $\frac{[0,+\infty]}{[0,+\infty]}$ cho

cheap but imprecise

Selective Flow-Sensitivity



FS : {x,y}







x	[I,+∞]
у	[Ⅰ,+∞]

FI : {z}



Selective Flow-Sensitivity



FS : {y,z}







у	[0,+∞]
z	[2,2]

fail to prove

 $FI: \{x\}$



Hard Search Problem

- Intractably large space, if not infinite
 - 2^{Var} different abstractions for FS
- Most of them are too imprecise or costly
 - $P(\{x,y,z\}) = \{ \emptyset, \{x\}, \{y\}, \{z\}, \{x,y\}, \{y,z\}, \{x,z\}, \{x,y,z\} \}$

Our Research

- How to automatically find a good abstraction?
 - pre-analysis approach [PLDI'14, TOPLAS'16]



• data-driven approaches [OOPSLA'15, SAS'16, APLAS'16]



Our Learning Approaches

- Learning via black-box optimization [OOPSLA'15]
- Learning via white-box optimization [APLAS'16]
- Learning from automatically labelled data [SAS'16]

. . .

• Learning with automatically generated features (in progress)

Static Analyzer

number of proved assertions

abstraction (e.g., a set of variables)

 $F(p, a) \Rightarrow n$

• Parameterized adaptation strategy

 $S_w: pgm \rightarrow 2^{Var}$

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• Learn a good parameter W from existing codebase

• Parameterized adaptation strategy

$$S_w: pgm \rightarrow 2^{Var}$$

• Learn a good parameter W from existing codebase

• For new program P, run static analysis with Sw(P)

I. Parameterized Strategy

$$S_w: pgm \rightarrow 2^{Var}$$

(I) Represent program variables as feature vectors.

(2) Compute the score of each variable.

(3) Choose the top-k variables based on the score.

(I) Features

• Predicates over variables:

$$f = \{f_1, f_2, \dots, f_5\}$$
 $(f_i: Var \rightarrow \{0, I\})$

- 45 simple syntactic features for variables: e.g,
 - local / global variable, passed to / returned from malloc, incremented by constants, etc

(I) Features

• Represent each variable as a feature vector:

 $f(x) = \langle f_1(x), f_2(x), f_3(x), f_4(x), f_5(x) \rangle$

 $f(x) = \langle 1, 0, 1, 0, 0 \rangle$ $f(y) = \langle 1, 0, 1, 0, 1 \rangle$ $f(z) = \langle 0, 0, 1, 1, 0 \rangle$

(2) Scoring

• The parameter w is a real-valued vector: e.g.,

$$w = \langle 0.9, 0.5, -0.6, 0.7, 0.3 \rangle$$

• Compute scores of variables:

score(x) = $\langle 1,0,1,0,0 \rangle \cdot \langle 0.9, 0.5, -0.6, 0.7, 0.3 \rangle = 0.3$ score(y) = $\langle 1,0,1,0,1 \rangle \cdot \langle 0.9, 0.5, -0.6, 0.7, 0.3 \rangle = 0.6$ score(z) = $\langle 0,0,1,1,0 \rangle \cdot \langle 0.9, 0.5, -0.6, 0.7, 0.3 \rangle = 0.1$

(3) Choose Top-k Variables

Choose the top-k variables based on their scores:
e.g., when k=2,



 In experiments, we chosen 10% of variables with highest scores.

2. Learn a Good Parameter



• Solve the optimization problem:

Find w that maximizes
$$\sum_{P_i} F(P_i, S_{\mathbf{w}}(P_i))$$

Learning via Random Sampling



Learning via Random Sampling



Bayesian Optimization

- A powerful method for solving difficult black-box optimization problems.
- Especially powerful when the objective function is expensive to evaluate.
- Key idea: use a probabilistic model to reduce the number of objective function evaluations.

Learning via Bayesian Optimization

repeat N times select a promising w using the model evaluate $\sum_{P_i} F(P_i, S_w(P_i))$ update the probabilistic model return best w found

- Probabilistic model: Gaussian processes
- Selection strategy: Expected improvement

Learning via Bayesian Optimization



Random Sampling vs Bayesian Optimization

Random sampling O Bayesian optimization



Effectiveness

- Implemented in Sparrow, an interval analyzer for C
- Evaluated on 30 open-source programs
 - 20 for training, 10 for testing

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Limitations

- While promising, the method has limitations:
 - black-box optimization is inherently inefficient
 - manual feature engineering is needed
- Follow-up work to overcome the limitations:
 - improving the efficiency [APLAS'16, SAS'16]
 - automating feature engineering [on-going]

Improving Efficiency

• A white-box optimization method [APLAS'16]

 $\mathcal{O}_P: \mathbb{J}_P \to \mathbb{R}.$

Find \mathbf{w}^* that minimizes $\sum_{j \in \mathbb{J}_P} (score_P^{\mathbf{w}}(j) - \mathcal{O}(j))^2$

• A supervised learning method [SAS'16]

	a	-a	b	-b	с	-c	i	-i
a	\star	Т	★	Т	Т	Т	\star	T
-a	T	\star	Т	\star	Т	Т	Т	T
b	\star	Т	★	Т	Т	Т	★	T
-b	T	★	Т	\star	Т	Т	Т	T
С	Τ	Т	Т	Т	\star	Т	Т	Т
-c	Τ	Т	Т	Т	Т	\star	Т	T
i	Т	Т	Т	Т	Т	Т	\star	Т
-i	T	\star	Т	\star	Т	Т	Т	\star

Manual Feature Engineering

Features

- The success of ML heavily depends on the "features"
- Feature engineering is nontrivial and time-consuming

leaf function

function containing malloc

 $\frac{1}{2}$

• Features do not generalize to other tasks

Type	#	Features
A	1	local variable
	2	global variable
	3	structure field
	4	location created by dynamic memory allocation
	5	defined at one program point
	6	location potentially generated in library code
	7	assigned a constant expression (e.g., $x = c1 + c2$)
	8	compared with a constant expression (e.g., $x < c$)
	9	compared with an other variable (e.g., $x < y$)
	10	negated in a conditional expression (e.g., if (!x))
	11	directly used in malloc (e.g., malloc(x))
	12	indirectly used in malloc (e.g., $y = x$; malloc(y))
	13	directly used in realloc (e.g., realloc(x))
	14	indirectly used in realloc (e.g., y = x; realloc(y))
	15	directly returned from malloc (e.g., x = malloc(e))
	16	indirectly returned from malloc
	17	directly returned from realloc (e.g., x = realloc(e))
	18	indirectly returned from realloc
	19	incremented by one (e.g., $x = x + 1$)
	20	incremented by a constant expr. (e.g., $x = x + (1+2)$)
	21	incremented by a variable (e.g., $x = x + y$)
	22	decremented by one (e.g., $x = x - 1$)
	23	decremented by a constant expr (e.g., $x = x - (1+2)$)
	24	decremented by a variable (e.g., $x = x - y$)
	25	multiplied by a constant (e.g., $x = x * 2$)
	26	multiplied by a variable (e.g., x = x * y)
	27	incremented pointer (e.g., p++)
	28	used as an array index (e.g., a[x])
	29	used in an array expr. (e.g., x[e])
	30	returned from an unknown library function
	31	modified inside a recursive function
	32	modified inside a local loop
	33	read inside a local loop
B	34	$1 \land 8 \land (11 \lor 12)$
	35	$2 \wedge 8 \wedge (11 \vee 12)$
	36	$1 \wedge (11 \vee 12) \wedge (19 \vee 20)$
	37	$2 \wedge (11 \vee 12) \wedge (19 \vee 20)$
	38	$1 \wedge (11 \vee 12) \wedge (15 \vee 16)$
	39	$2 \wedge (11 \lor 12) \wedge (15 \lor 16)$
	40	$(11 \lor 12) \land 29$ $(15 \lor 16) \land 29$
	41	$(15 \lor 10) \land 29$
	42	$1 \land (19 \lor 20) \land 33$
	43	$2 \land (19 \lor 20) \land 33$
	44	$1 \land (19 \lor 20) \land \neg 33$
	45	$2 \wedge (19 \vee 20) \wedge \neg 33$

3 function containing realloc function containing a loop function containing an if statement function containing a switch statement function using a string-related library function write to a global variable read a global variable 9 10 write to a structure field 11 read from a structure field directly return a constant expression 1213 indirectly return a constant expression 14 directly return an allocated memory 15 indirectly return an allocated memory 16 directly return a reallocated memory 17 indirectly return a reallocated memory 18 return expression involves field access return value depends on a structure field 19 20return void directly invoked with a constant 2122 constant is passed to an argument 23invoked with an unknown value 24 functions having no arguments 25 functions having one argument 26functions having more than one argument 27functions having an integer argument 28functions having a pointer argument 29 | functions having a structure as an argument 30 B $2 \land (21 \lor 22) \land (14 \lor 15)$ $31 \mid 2 \land (21 \lor 22) \land \neg (14 \lor 15)$ $32 \mid 2 \land 23 \land (14 \lor 15)$ **33** $2 \wedge 23 \wedge \neg (14 \vee 15)$ 34 $2 \wedge (21 \vee 22) \wedge (16 \vee 17)$ 35 $2 \wedge (21 \vee 22) \wedge \neg (16 \vee 17)$ 36 $2 \wedge 23 \wedge (16 \vee 17)$ 37 $2 \wedge 23 \wedge \neg (16 \vee 17)$ 38 | $(21 \lor 22) \land \neg 23$

-		
Type	#	Features
A	1	used in array declarations (e.g., a[c])
	2	used in memory allocation (e.g., malloc(c))
	3	used in the righthand-side of an assignment (e.g., $x = c$)
	4	used with the less-than operator (e.g, $x < c$)
	5	used with the greater-than operator (e.g., $x > c$)
	6	used with \leq (e.g., x \leq c)
	7	used with \geq (e.g., $x \geq c$)
	8	used with the equality operator (e.g., $x == c$)
	9	used with the not-equality operator (e.g., $x ! = c$)
	10	used within other conditional expressions (e.g., x < c+y)
	11	used inside loops
	12	used in return statements (e.g., return c)
	13	constant zero
В	14	$(1 \lor 2) \land 3$
	15	$(1 \lor 2) \land (4 \lor 5 \lor 6 \lor 7)$
	16	$(1 \lor 2) \land (8 \lor 9)$
	17	$(1 \lor 2) \land 11$
	18	$(1 \lor 2) \land 12$
	19	$13 \wedge 3$
	20	$13 \land (4 \lor 5 \lor 6 \lor 7)$
	21	$13 \land (8 \lor 9)$
	22	$13 \land 11$
	23	$13 \land 12$

flow-sensitivity

context-sensitivity

widening thresholds

Automatic Feature Generation

Before [OOPSLA'15,SAS'16,APLAS'16]



New method



Partial Flow-Sensitive Analysis

- A query-based, partially flow-sensitive interval analysis
- The analysis uses a query-classifier $C : Query \rightarrow \{1,0\}$

- x = 0; y = 0; z = input(); w = 0; y = x; y++;
- 3 assert (y > 0); // Query 1
- 4 assert (z > 0); // Query 2
- 5 | assert (w == 0); // Query 3

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x = 0; y = 0; z = input(); w = 0;y = x; y++;|assert (y > 0); // Query 1 provable assert (z > 0); // Query 2 unprovable assert (w == 0); // Query 3 unprovable

Partial Flow-Sensitive Analysis

- A query-based, partially flow-sensitive interval analysis
- The analysis uses a query-classifier $C : Query \rightarrow \{1,0\}$

x = 0;	y = 0; z =	<pre>input();</pre>	W	= 0;
y = x;	y++;			
assert	(y > 0);	// Query	1	provable
assert	(z > 0);	// Query	2	unprovable
assert	(w == 0);	// Query	3	unprovable

	flow-sensitive result	flow-insensitive result
line	abstract state	abstract state
1	$ \{x \mapsto [0,0], y \mapsto [0,0]\} $	
2	$\{x \mapsto [0,0], y \mapsto [1,1]\}$	
3	$\{x \mapsto [0,0], y \mapsto [1,1]\}$	$\{z\mapsto [0,0], w\mapsto [0,0]\}$
4	$\{x \mapsto [0,0], y \mapsto [1,1]\}$	
5	$ \{x \mapsto [0,0], y \mapsto [1,1]\} $	

Standard binary classification:

 $\{(q_i, b_i)\}_{i=1}^n$

Standard binary classification:



Standard binary classification:



Standard binary classification:



- Success relies on how we convert queries to feature vectors
- This feature engineering has been done manually

Conversion from Queries to Feature Vectors

- A set of feature features $\Pi = \{\pi_1, \ldots, \pi_k\}$
 - a feature encodes a property about queries
- A procedure to check whether a query satisfies a feature

 $\mathsf{match}: \mathit{Query} \times \mathit{Feature} \to \mathbb{B}$

• The feature vector of a query q:

 $\langle \mathsf{match}(q,\pi_1),\ldots,\mathsf{match}(q,\pi_k) \rangle$

Automatic Feature Generation

- Generate feature programs by running reducer
 - small pieces of code that minimally describe when it is worth increasing the precision
- Represent them by abstract data-flow graphs
 - generalized form of feature programs

Generating Feature Programs

- By running a program reducer: e.g., C-Reduce [PLDI'12] reduce : $\mathbb{P} \times (\mathbb{P} \to \mathbb{B}) \to \mathbb{P}$
- Feature-preserving condition: $\phi(P) \equiv FI(P) = unproven \land FS(P) = proven$

Generalize to Abstract Data-Flow Graphs



- The right level of abstraction depends on an analysis
- We choose the best abstraction using a combination of searching and cross-validation

Feature Generation

• Apply the method on codebases:



$$\Rightarrow \Pi = \{\pi_1, \ldots, \pi_k\}$$

Codebase

Matching Algorithm

 $\mathsf{match}: \mathit{Query} \times \mathit{Feature} \to \mathbb{B}$



Matching Algorithm

 $\mathsf{match}: \mathit{Query} \times \mathit{Feature} \to \mathbb{B}$



Subgraph inclusion:

 $(N_1, E_1) \subseteq (N_2, E_2) \iff N_1 \subseteq N_2 \land E_1 \subseteq E_2^*$

$$\{(v_i, b_i)\}_{i=1}^n$$
$$\bigcup$$
$$\mathcal{C} : \mathbb{B}^k \to \mathbb{B}$$

 $= \{\pi_1, \ldots, \pi_k\}$

Experiments

Effectiveness of partially flow-sensitive analysis

	Query Pro	ediction		Analysis								Comparison			
				Prove			Sec					Oh et al	. [38]		
Trial	Precision	Recall	Fli	FSi	Ours	Fli	FSi	Ours	Quality	Cost	Self	Quality	Cost		
1	92.6 %	77.9 %	5,340	6,053	5,973	38.2	564.0	55.3	88.7 %	1.4x	88.7 %	85.2%	1.5x		
2	78.8 %	73.3 %	2,972	3,373	3,262	16.3	460.5	25.7	72.3 %	1.5x	72.0 %	41.6%	1.9x		
3	66.7 %	73.3 %	3,984	4,668	4,559	27.3	1,635.6	176.2	84.0 %	6.4x	82.7 %	89.9%	3.2x		
4	88.7 %	68.8 %	4,600	5,450	5,307	38.1	688.2	59.6	83.1 %	1.5x	83.5 %	60.7%	1.9x		
5	89.9 %	79.4 %	2,517	2,971	2,945	10.9	325.9	18.9	94.2 %	1.7x	94.0 %	47.8%	2.1x		
TOTAL	81.5 %	73.9 %	19,413	22,515	22,046	131.1	3,674.4	336.0	84.8 %	2.5x	84.6 %	68.4%	2.1x		

Effectiveness of partially relational analysis

	Query Prediction					Anal	lysis	Comparison					
				Prove		Sec						Heo et al	. [21]
Trial	Precision	Recall	FSi	IMPCT	Ours	FSi	IMPCT	Ours	Quality	Cost	Self	Quality	Cost
1	74.8 %	81.3 %	3,678	3,806	3,789	140.7	389.8	189.5	86.7 %	1.3 x	54.2 %	100.0 %	3.0 x
2	84.1 %	82.6 %	5,845	6,004	5,977	613.5	18,022.9	775.5	83.0 %	1.3 x	65.5 %	30.2 %	0.9 x
3	82.8 %	73.0 %	1,926	2,079	2,036	315.2	2,396.9	460.2	71.9 %	1.5 x	95.7 %	92.2 %	1.1 x
4	77.6 %	85.2 %	2,221	2,335	2,313	72.7	495.1	141.2	80.7 %	1.9 x	67.2 %	100.0 %	2.0 x
5	71.6 %	78.4 %	2,886	2,962	2,946	148.9	557.2	210.2	78.9 %	1.4 x	59.9 %	96.1 %	2.3 x
TOTAL	79.0 %	79.9 %	16,556	17,186	17,061	1,291.0	21,861.9	1,776.6	80.2 %	1.4 x	67.7 %	80.0 %	1.4 x

Summary

- Choosing a good abstraction is a key challenge in static program analysis
- New data-driven approach is promising
- Further information:

http://prl.korea.ac.kr

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Thank you