# Static Analysis-based Repair of Memory Errors in C Programs

Hakjoo Oh Korea University

2/25/2025@IFIP WG 2.4 Meeting 70 (Singapore)

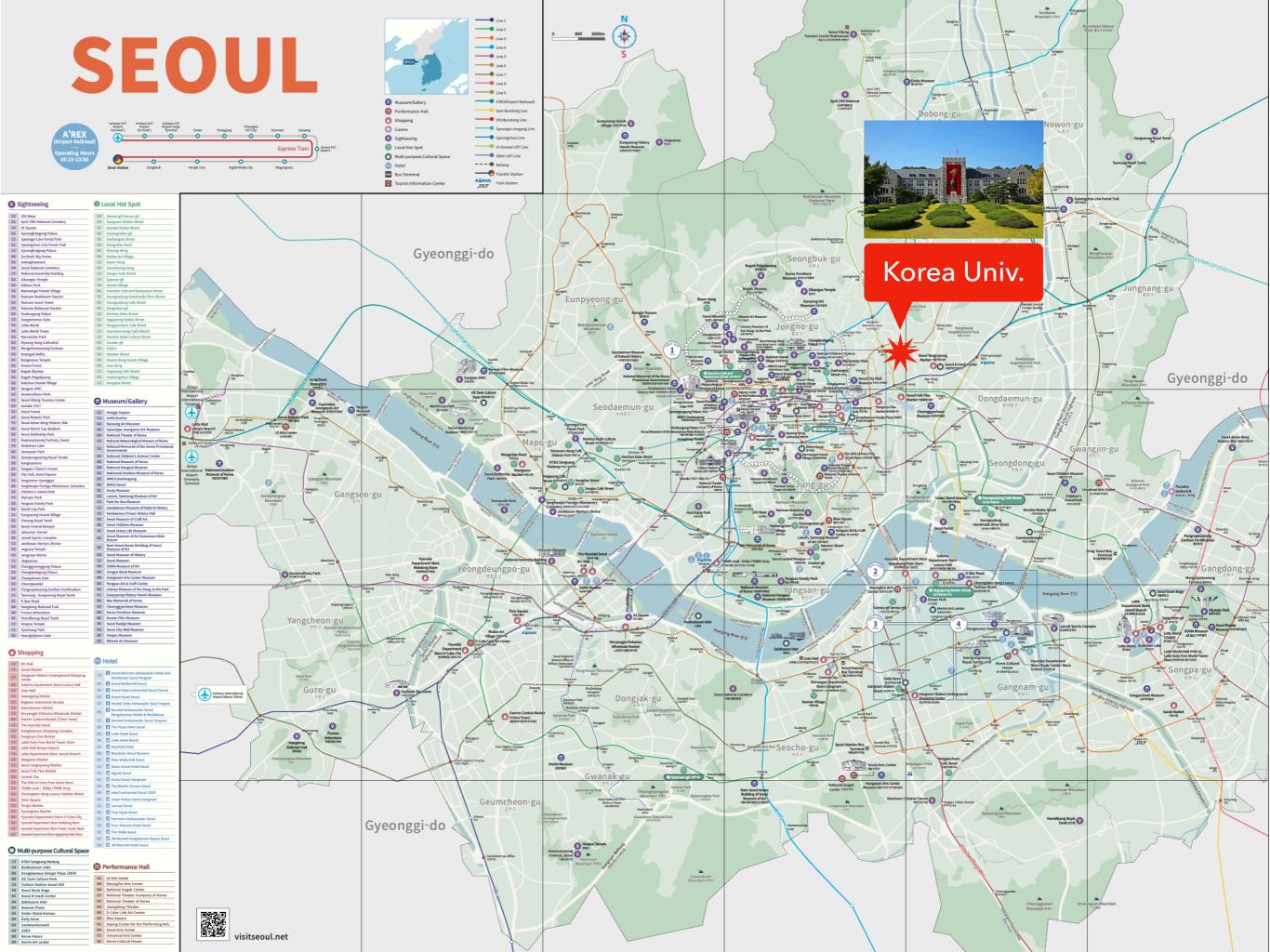
### PL/SE Research @Korea Univ.

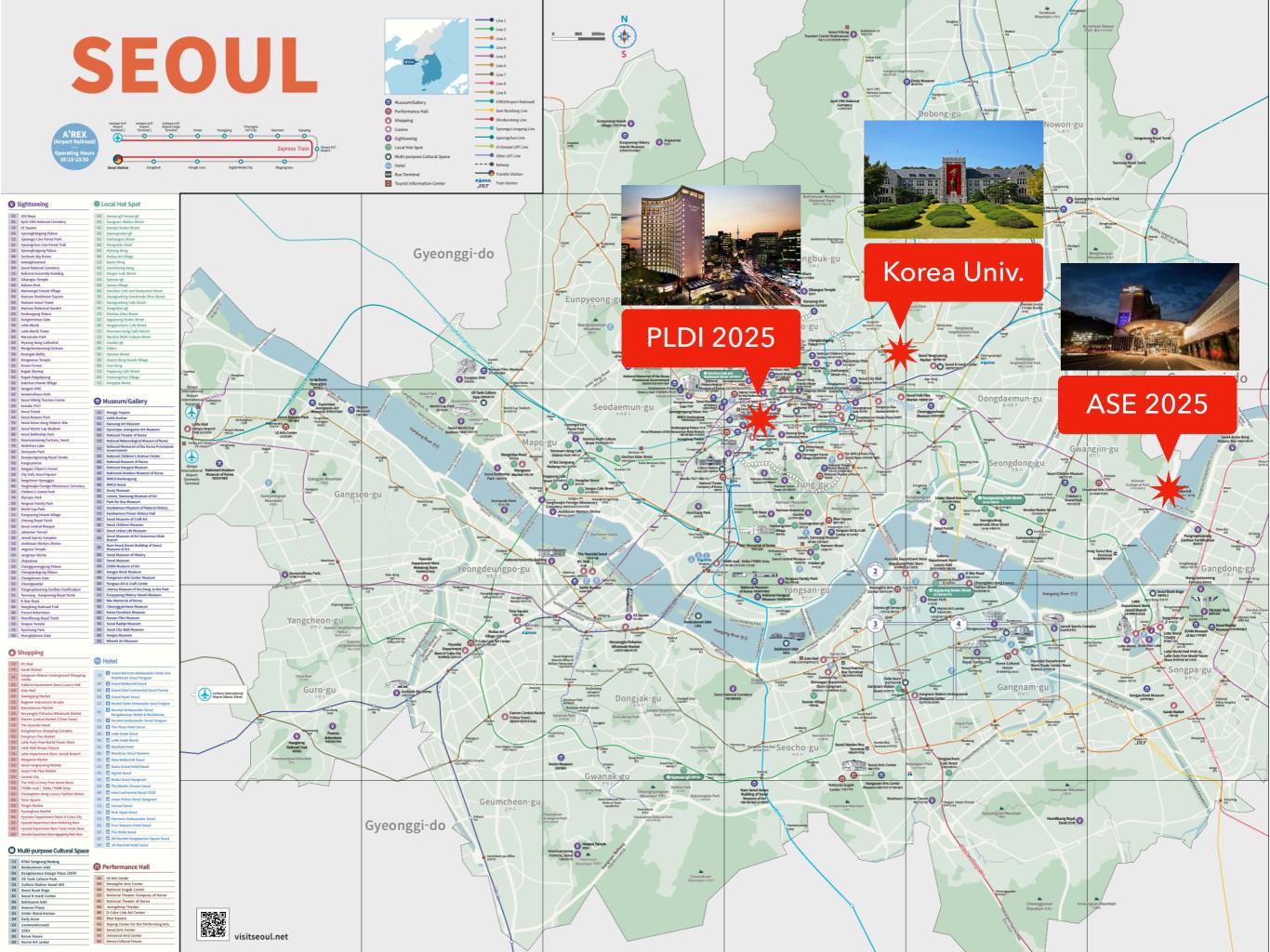
- Members: ~15 graduate students
- Research area: intersection of programming languages
   (PL) and software engineering (SE)
  - program analysis and testing
  - program synthesis and repair
- Publication: PL, SE, and Security



- PL: POPL('22),PLDI('12,'14,'20,'24),OOPSLA('15,'17a,'17b,'18a,'18b,'19,'20,'23,'24a,'24b,25)
- SE: ICSE('17,'18,'19,'20,'21'22a,'22b,'23a,'23b,'23c), FSE('18,'19,'20,'21,'22,'23), ASE('18,'24a,'24b)
- Security: IEEE S&P('17,'20), USENIX Security('21,'23)

#### http://kupl.github.io





#### My IFIP Talks

- Meeting 67, York Harbor (April 23-27, 2023)
- Meeting 69, Lugano (May 12-16, 2024)

#### **Data-Driven Static Analysis**

Hakjoo Oh



25 April 2023 @IFIP WG 2.4 Meeting 67, York Harbor

PL4XGL: A Programming Language Approach to Explainable Graph Learning

> Hakjoo Oh Korea University

(co-work with <u>Minseok Jeon</u> and Jihyeok Park)

IFIP WG 2.4 Meeting @Lugano, Switzerland

#### My IFIP Talks

- Meeting 67, York Harbor (April 23-27, 2023)
- Meeting 69, Lugano (May 12-16, 2024)

#### **Data-Driven Static Analysis**

Hakjoo Oh



25 April 2023 @IFIP WG 2.4 Meeting 67, York Harbor

PL4XGL: A Programming Language Approach to Explainable Graph Learning

> Hakjoo Oh Korea University

(co-work with <u>Minseok Jeon</u> and Jihyeok Park)

IFIP WG 2.4 Meeting @Lugano, Switzerland

Today: automated program repair (APR)

- Automatic Diagnosis and Correction of Logical Errors for Functional Programming Assignments. OOPSLA 2018
- Context-Aware and Data-Driven Feedback Generation for Programming Assignments. ESEC/FSE 2021
- MemFix: Static Analysis-Based Repair of Memory Deallocation Errors for C. ESEC/FSE 2018
- SAVER: Scalable, Precise, and Safe Memory-Error Repair. ICSE 2020 (deployed in industry)
- NPEX: Repairing Java Null Pointer Exceptions without Tests. ICSE 2022
- PyTER: Effective Program Repair for Python Type Errors. ESEC/FSE 2022
- SmartFix: Fixing Vulnerable Smart Contracts by Accelerating Generate-and-Verify Repair using Statistical Models. ESEC/FSE 2023
- Reducing the Cost of LLM-based APR via Execution-Guided Static Analysis. In progress
- Accurate Detection of Overfitting Patches in Automated Program Repair through Semantic Anti-Patterns. In progress

- Automatic Diagnosis and Correction of Logical Errors for Functional Programming Assignments. OOPSLA 2018
- Context-Aware and Data-Driven Feedback Generation for Programming Assignments. ESEC/FSE 2021
- MemFix: Static Analysis-Based Repair of Memory Deallocation Errors for C. ESEC/FSE 2018
- SAVER: Scalable, Precise, and Safe Memory-Error Repair. ICSE 2020 (deployed in industry)
- NPEX: Repairing Java Null Pointer Exceptions without Tests. ICSE 2022
- PyTER: Effective Program Repair for Python Type Errors. ESEC/FSE 2022
- SmartFix: Fixing Vulnerable Smart Contracts by Accelerating Generate-and-Verify Repair using Statistical Models. ESEC/FSE 2023
- Reducing the Cost of LLM-based APR via Execution-Guided Static Analysis. In progress
- Accurate Detection of Overfitting Patches in Automated Program Repair through Semantic Anti-Patterns. In progress

- Automatic Diagnosis and Correction of Logical Errors for Functional Programming Assignments. OOPSLA 2018
- Context-Aware and Data-Driven Feedback Generation for Programming Assignments. ESEC/FSE 2021
- MemFix: Static Analysis-Based Repair of Memory Deallocation Errors for C. ESEC/FSE 2018
- SAVER: Scalable, Precise, and Safe Memory-Error Repair. ICSE 2020 (deployed in industry)
- NPEX: Repairing Java Null Pointer Exceptions without Tests. ICSE 2022
- PyTER: Effective Program Repair for Python Type Errors. ESEC/FSE 2022
- SmartFix: Fixing Vulnerable Smart Contracts by Accelerating Generate-and-Verify Repair using Statistical Models. ESEC/FSE 2023
- Reducing the Cost of LLM-based APR via Execution-Guided Static Analysis. In progress
- Accurate Detection of Overfitting Patches in Automated Program Repair through Semantic Anti-Patterns. In progress

- Automatic Diagnosis and Correction of Logical Errors for Functional Programming Assignments. OOPSLA 2018
- Context-Aware and Data-Driven Feedback Generation for Programming Assignments. ESEC/FSE 2021
- MemFix: Static Analysis-Based Repair of Memory Deallocation Errors for C. ESEC/FSE 2018
- SAVER: Scalable, Precise, and Safe Memory-Error Repair. ICSE 2020 (deployed in industry)
- NPEX: Repairing Java Null Pointer Exceptions without Tests. ICSE 2022
- PyTER: Effective Program Repair for Python Type Errors. ESEC/FSE 2022
- SmartFix: Fixing Vulnerable Smart Contracts by Accelerating Generate-and-Verify Repair using Statistical Models. ESEC/FSE 2023
- Reducing the Cost of LLM-based APR via Execution-Guided Static Analysis. In progress
- Accurate Detection of Overfitting Patches in Automated Program Repair through Semantic Anti-Patterns. In progress

- Automatic Diagnosis and Correction of Logical Errors for Functional Programming Assignments. OOPSLA 2018
- Context-Aware and Data-Driven Feedback Generation for Programming Assignments. ESEC/FSE 2021
- MemFix: Static Analysis-Based Repair of Memory Deallocation Errors for C. ESEC/FSE 2018
- SAVER: Scalable, Precise, and Safe Memory-Error Repair. ICSE 2020 (deployed in industry)
- NPEX: Repairing Java Null Pointer Exceptions without Tests. ICSE 2022
- PyTER: Effective Program Repair for Python Type Errors. ESEC/FSE 2022
- SmartFix: Fixing Vulnerable Smart Contracts by Accelerating Generate-and-Verify Repair using Statistical Models. ESEC/FSE 2023
- Reducing the Cost of LLM-based APR via Execution-Guided Static Analysis. In progress
- Accurate Detection of Overfitting Patches in Automated Program Repair through Semantic Anti-Patterns. In progress

- Automatic Diagnosis and Correction of Logical Errors for Functional Programming Assignments. OOPSLA 2018
- Context-Aware and Data-Driven Feedback Generation for Programming Assignments. ESEC/FSE 2021
- MemFix: Static Analysis-Based Repair of Memory Deallocation Errors for C. ESEC/FSE 2018
- SAVER: Scalable, Precise, and Safe Memory-Error Repair. ICSE 2020 (deployed in industry)
- NPEX: Repairing Java Null Pointer Exceptions without Tests. ICSE 2022
- PyTER: Effective Program Repair for Python Type Errors. ESEC/FSE 2022
- SmartFix: Fixing Vulnerable Smart Contracts by Accelerating Generate-and-Verify Repair using Statistical Models. ESEC/FSE 2023
- Reducing the Cost of LLM-based APR via Execution-Guided Static Analysis. In progress
- Accurate Detection of Overfitting Patches in Automated Program Repair through Semantic Anti-Patterns. In progress

- Automatic Diagnosis and Correction of Logical Errors for Functional Programming Assignments. OOPSLA 2018
- Context-Aware and Data-Driven Feedback Generation for Programming Assignments. ESEC/FSE 2021
- MemFix: Static Analysis-Based Repair of Memory Deallocation Errors for C. ESEC/FSE 2018
- SAVER: Scalable, Precise, and Safe Memory-Error Repair. ICSE 2020 (deployed in industry)
- NPEX: Repairing Java Null Pointer Exceptions without Tests. ICSE 2022
- PyTER: Effective Program Repair for Python Type Errors. ESEC/FSE 2022
- SmartFix: Fixing Vulnerable Smart Contracts by Accelerating Generate-and-Verify Repair using Statistical Models. ESEC/FSE 2023
- Reducing the Cost of LLM-based APR via Execution-Guided Static Analysis. In progress
- Accurate Detection of Overfitting Patches in Automated Program Repair through Semantic Anti-Patterns. In progress

Our approach: Static analysis-based program repair

- Automatic Diagnosis and Correction of Logical Errors for Functional Programming Assignments. OOPSLA 2018
- Context-Aware and Data-Driven Feedback Generation for Programming Assignments. ESEC/FSE 2021
- MemFix: Static Analysis-Based Repair of Memory Deallocation Errors for C. ESEC/FSE 2018
- SAVER: Scalable, Precise, and Safe Memory-Error Repair. ICSE 2020 (deployed in industry)
- NPEX: Repairing Java Null Pointer Exceptions without Tests. ICSE 2022
- PyTER: Effective Program Repair for Python Type Errors. ESEC/FSE 2022
- SmartFix: Fixing Vulnerable Smart Contracts by Accelerating Generate-and-Verify Repair using Statistical Models. ESEC/FSE 2023
- Reducing the Cost of LLM-based APR via Execution-Guided Static Analysis. In progress
- Accurate Detection of Overfitting Patches in Automated Program Repair through Semantic Anti-Patterns. In progress

Our approach: Static analysis-based program repair

## Memory Errors in C Programs

 Memory-leak (ML), use-after-free (UAF), and double-free (DF) are prevalent in real-world C programs

Repository	#commits	ML	DF	UAF	Total	*-overflow
linux	721,119	3,740	821	1,986	6,363	5,092
openssl	21,009	220	36	12	264	61
numpy	17,008	58	2	2	59	53
php	105,613	1,129	148	197	1,449	649
git	49,475	350	19	95	442	258

## Memory Errors in C Programs

 Memory-leak (ML), use-after-free (UAF), and double-free (DF) are prevalent in real-world C programs

Repository	#commits	ML	DF	UAF	Total	*-overflow
linux	721,119	3,740	821	1,986	6,363	5,092
openssl	21,009	220	36	12	264	61
numpy	17,008	58	2	2	59	53
php	105,613	1,129	148	197	1,449	649
git	49,475	350	19	95	442	258

• Significant sources of security vulnerabilities

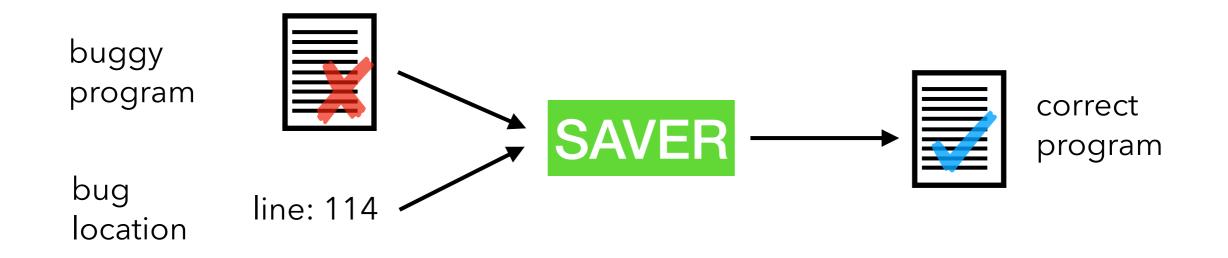
CVE-2017-9798 Optionsbleed - Apache memory leak

Linux kernel: CVE-2017-6074: DCCP double-free vulnerability (local root)					
From: Andrey Konovalov <andreyknvl ()="" com="" google=""> Date: Wed, 22 Feb 2017 14:28:35 +0100</andreyknvl>					
ні,					
This is an announcement about CVE-2017-6074 [1] which is a double-free vulnerability I found in the Linux kernel. It can be exploited to gain kernel code execution from an unprivileged processes.					
	- 1				

/ulnerability Details : <u>CVE-2017-11274</u>			
Adobe Digital Editions 4.5.4 and earlier has an earlier has an earlier has an earlier has an earlier bate : 2017-08-11 Last Update Date : 2017	•	after free vulne	erability.
ollapse All Expand All Select Select&Copy earch Twitter Search YouTube Search Google	Scroll To	- Comments	▼ Exter
- CVSS Scores & Vulnerability Types			
CVSS Score 10.0			

#### Goal

- Long-term goal: Fully automated detection and repair
- This talk: SAVER, a system to automatically fix memory errors



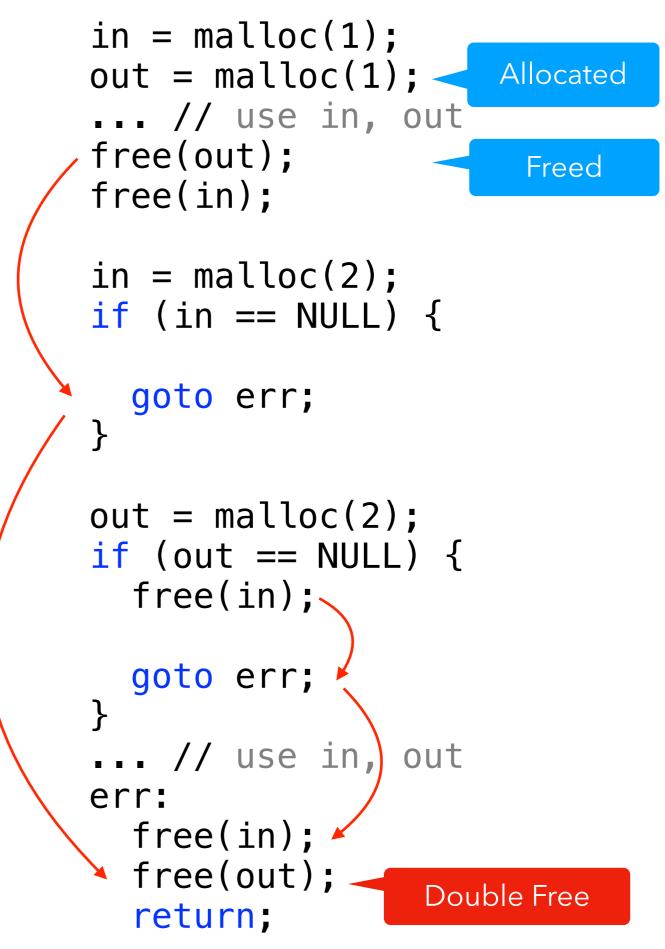
- To be practical, SAVER is designed to be
  - Scalable: Capable of handling large, industry programs
  - **Precise**: Effectively fixes diverse bugs with high fix rates
  - Safe: Generated patches are safe, not introducing new errors

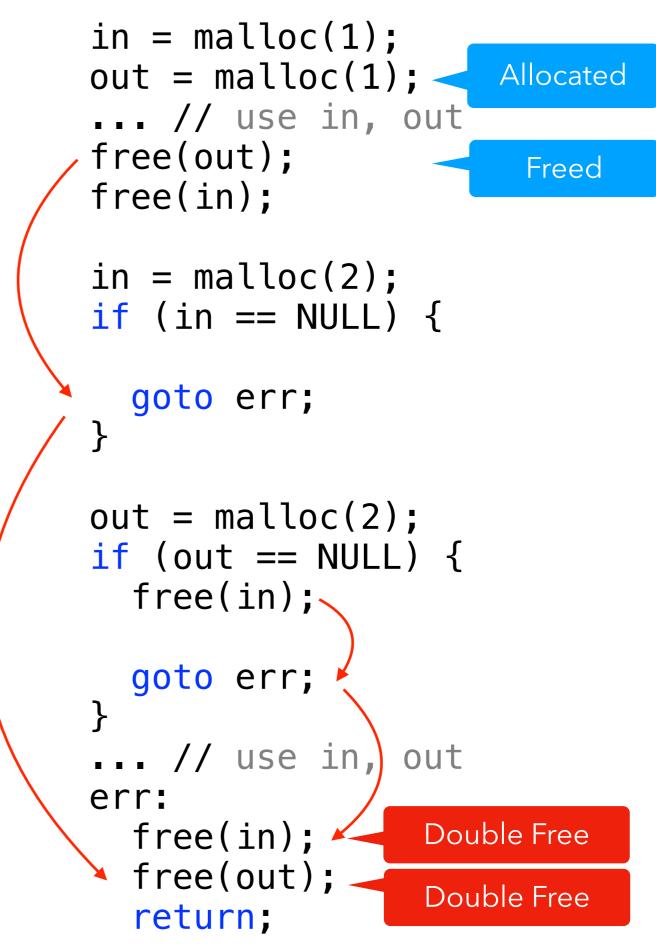
```
in = malloc(1);
out = malloc(1);
... // use in, out
free(out);
free(in);
in = malloc(2);
if (in == NULL) {
  goto err;
}
out = malloc(2);
if (out == NULL) {
  free(in);
  goto err;
}
... // use in, out
err:
  free(in);
  free(out);
  return;
```

```
in = malloc(1);
                    Allocated
out = malloc(1);
... // use in, out
free(out);
free(in);
in = malloc(2);
if (in == NULL) {
  goto err;
}
out = malloc(2);
if (out == NULL) {
  free(in);
  goto err;
}
... // use in, out
err:
  free(in);
  free(out);
  return;
```

```
in = malloc(1);
                     Allocated
out = malloc(1);
... // use in, out
free(out);
                      Freed
free(in);
in = malloc(2);
if (in == NULL) {
  goto err;
}
out = malloc(2);
if (out == NULL) {
  free(in);
  goto err;
}
... // use in, out
err:
  free(in);
  free(out);
  return;
```

```
in = malloc(1);
                     Allocated
out = malloc(1);
... // use in, out
free(out);
                       Freed
free(in);
in = malloc(2);
if (in == NULL) {
  goto err;
}
out = malloc(2);
if (out == NULL) {
  free(in);
  goto err;
}
... // use in, out
err:
  free(in);
  free(out);
                  Double Free
  return;
```





#### USB: fix double frees in error code paths of ipaq driver

the error code paths can be enter with buffers to freed buffers. Serial core would do a kfree() on memory already freed.

Signed-off-by: Oliver Neukum <oneukum@suse.de>
Signed-off-by: Greg Kroah-Hartman <gregkh@suse.de>

master > v4.15-rc1 ... v2.6.24-rc1

Oliver Neukum committed with gregkh on 18 Sep 2007

```
in = malloc(1);
out = malloc(1);
... // use in, out
free(out);
free(in);
```

```
in = malloc(2);
if (in == NULL) {
    out = NULL;
    goto err;
}
```

```
out = malloc(2);
if (out == NULL) {
   free(in);
   in = NULL;
   goto err;
}
... // use in, out
err:
   free(in);
   free(out);
   return;
```

1 par

#### USB: fix double frees in error code paths of ipaq driver

the error code paths can be enter with buffers to freed buffers. Serial core would do a kfree() on memory already freed.

Signed-off-by: Oliver Neukum <oneukum@suse.de>
Signed-off-by: Greg Kroah-Hartman <gregkh@suse.de>

master > v4.15-rc1 ... v2.6.24-rc1

Oliver Neukum committed with **gregkh** on 18 Sep 2007

# Challenge 1: Difficult to ensure that bugs are fixed correctly

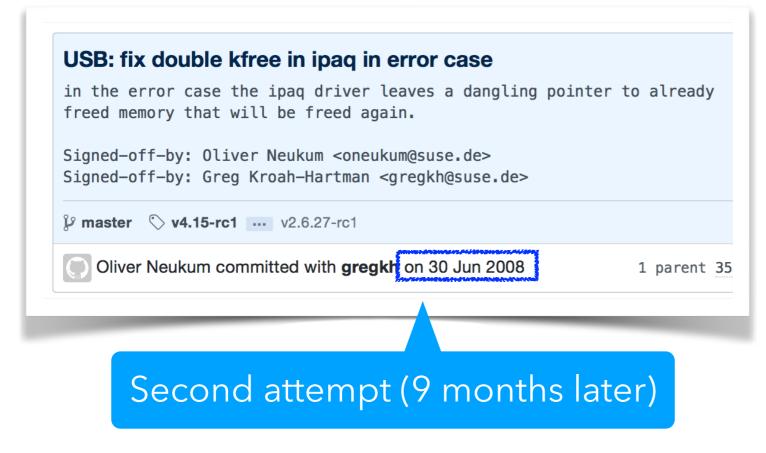
```
in = malloc(1);
out = malloc(1);
... // use in, out
free(out);
free(in);
```

```
in = malloc(2);
if (in == NULL) {
    out = NULL;
    goto err;
}
```

```
out = malloc(2);
if (out == NULL) {
   free(in);
   in = NULL;
   goto err;
}
... // use in, out
err:
   free(in);
   free(out);
```

return;

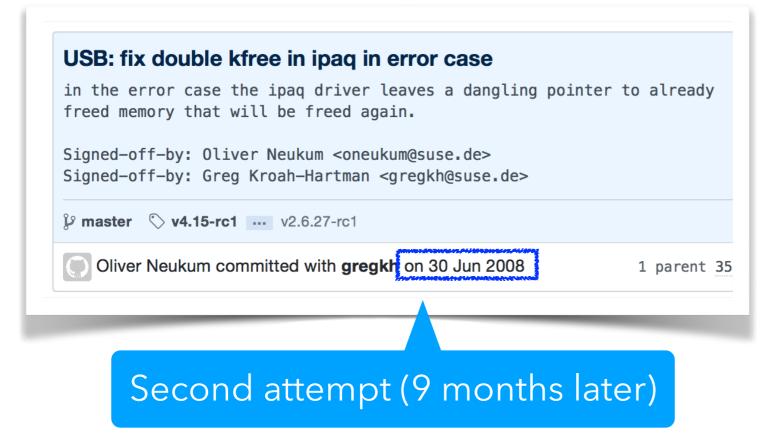
1 par



```
in = malloc(1);
out = malloc(1);
... // use in, out
free(out);
free(in);
in = malloc(2);
if (in == NULL) {
  out = NULL;
  goto err;
}
free(out);
out = malloc(2);
if (out == NULL) {
  free(in);
  in = NULL;
  goto err;
}
... // use in, out
err:
  free(in);
  free(out);
  return;
```

memory leak

# Challenge 2: Patches may introduce new errors



in = malloc(1);out = malloc(1); ... // use in, out free(out); free(in); in = malloc(2);if (in == NULL) { out = NULL; goto err; } free(out); out = malloc(2); if (out == NULL) { free(in); in = NULL; goto err; } ... // use in, out err: free(in); free(out); return;

fix <sup>·</sup>	for	a	memory	leak	in an	error	case	introduced	by	fix	for	doub	le 1	free
------------------	-----	---	--------	------	-------	-------	------	------------	----	-----	-----	------	------	------

The fix NULLed a pointer without freeing it.

Signed-off-by: Oliver Neukum <oneukum@suse.de>
Reported-by: Juha Motorsportcom <juha\_motorsportcom@luukku.com>
Signed-off-by: Linus Torvalds <torvalds@linux-foundation.org>

master > v4.15-rc1 --- v2.6.27-rc1

Oliver Neukum committed with torvalds on 27 Jul 2008

1 parent 9ee08c2

#### Third attempt (10 months after bug detection)

```
in = malloc(1);
out = malloc(1);
... // use in, out
free(out);
free(in);
out = NULL;
in = malloc(2);
if (in == NULL) {
  out = NULL;
  goto err;
}
free(out);
out = malloc(2);
if (out == NULL) {
  free(in);
  in = NULL;
  goto err;
}
... // use in, out
err:
  free(in);
  free(out);
  return;
```

#### Challenge 3: The resulting patches are often not of high quality

Ine	fix NULLed a pointer without freeing it.
Repo	ed—off—by: Oliver Neukum <oneukum@suse.de> rted—by: Juha Motorsportcom <juha_motorsportcom@luukku.com> ed—off—by: Linus Torvalds <torvalds@linux—foundation.org></torvalds@linux—foundation.org></juha_motorsportcom@luukku.com></oneukum@suse.de>
ې ma	aster 🟷 v4.15-rc1 v2.6.27-rc1
$\bigcirc$	Oliver Neukum committed with torvalds on 27 Jul 2008 1 parent 9ee08c2
	Third attempt
	(10 months after bug detection)

in = malloc(1);out = malloc(1); ... // use in, out free(out); free(in); out = NULL; in = malloc(2);if (in == NULL) { out = NULL; goto err; } free(out); out = malloc(2); if (out == NULL) { free(in); in = NULL; goto err; } ... // use in, out err: free(in); free(out); return;

#### SAVER-Generated Patch

```
in = malloc(1);
                                                   in = malloc(1);
out = malloc(1);
                                                   out = malloc(1);
... // use in, out
                                                   ... // use in, out
free(out);
                                                   free(out);
free(in);
                                                   free(in);
in = malloc(2);
                                                   in = malloc(2);
                                                   if (in == NULL) {
if (in == NULL) {
                             SAVER
 goto err;
                                                     goto err;
}
                                                   }
                                                   free(out);
                          ✓ Fast (few mins)
out = malloc(2);
                                                   out = malloc(2);
if (out == NULL) {
                          ✓ Safety guarantee
                                                   if (out == NULL) {
                                                   free(in);
  free(in);
                                                     goto err;
  goto err;
}
... // use in, out
                                                   ... // use in, out
err:
                                                   err:
  free(in); // double-free
                                                     free(in);
  free(out);// double-free
                                                     free(out);
  return;
                                                     return;
```

```
int append_data (Node *node, int *ndata) {
1
       if (!(Node *n = malloc(sizeof(Node)))
2
           return -1; // failed to be appended
3
       n->data = ndata;
4
       n->next = node->next; node->next = n;
5
       return 0; // successfully appended
6
   }
7
8
   Node *lx = ... // a linked list
9
   Node *ly = ... // a linked list
10
   for (Node *node = 1x; node != NULL; node = node->next) {
11
       int *dptr = malloc(sizeof(int));
12
       if (!dptr) return;
13
       *dptr = *(node->data);
14
   (-) append_data(ly, dptr); // potential memory-leak
15
16
    }
17
                     Memory Leak:
                     An object allocated at line 12
           Infer
```

becomes unreachable after line 15

```
int append_data (Node *node, int *ndata) {
1
       if (!(Node *n = malloc(sizeof(Node)))
2
           return -1; // failed to be appended
3
       n->data = ndata;
4
       n \rightarrow next = node \rightarrow next; node \rightarrow next = n;
5
       return 0; // successfully appended
6
    }
7
                                               Normal execution
8
   Node *lx = ... // a linked list
9
   Node *ly = ... // a linked list
10
    for (Node *node = 1x; node != NULL, node = node->next) {
11
       int *dptr = malloc(sizeof(int));
12
       if (!dptr) return;
13
       *dptr = *(node->data);
14
    (-) append_data(ly, dptr); // potential memory-leak
15
16
    }
17
                     Memory Leak:
                      An object allocated at line 12
           Infer
                      becomes unreachable after line 15
```

```
int append_data (Node *node, int *ndata) {
1
       if (!(Node *n = malloc(sizeof(Nod)))
2
           return -1; #7 failed to be appended
3
       n->data = ndata;
4
       n \rightarrow next = node \rightarrow next; node \rightarrow next = r;
5
       return 0; // successfully appended
6
    }
7
                                               Buggy execution
8
   Node *lx = ... // a linked list
9
   Node *ly = ... // a linked list
10
    for (Node *node = 1x; node != NULL, node = node->next) {
11
       int *dptr = malloc(sizeof(int));
12
       if (!dptr) return;
13
       *dptr = *(node->data);
14
    (-) append_data(ly, dptr); // potential memory-leak
15
16
    }
17
                     Memory Leak:
                      An object allocated at line 12
           Infer
                      becomes unreachable after line 15
```

```
int append_data (Node *node, int *ndata) {
1
       if (!(Node *n = malloc(sizeof(Node)))
2
           return -1; // failed to be appended
3
       n->data = ndata;
4
       n->next = node->next; node->next = n;
5
       return 0; // successfully appended
6
   }
7
8
   Node *lx = ... // a linked list
9
   Node *ly = ... // a linked list
10
   for (Node *node = 1x; node != NULL; node = node->next) {
11
       int *dptr = malloc(sizeof(int));
12
       if (!dptr) return;
13
       *dptr = *(node->data);
14
   (-) append_data(ly, dptr); // potential memory-leak
15
   (+) if ((append_data(ly, dptr)) == -1) free(dptr);
16
    }
17
                     Memory Leak:
```

Infer

An object allocated at <u>line 12</u> becomes unreachable after <u>line 15</u>

## cf) SAVER vs. Generative Al

• LLMs do not guarantee safety, e.g., GPT4-generated patch:

```
int append_data (Node *node, int *ndata) {
1
       if (!(Node *n = malloc(sizeof(Node)))
2
           return -1; // failed to be appended
3
       n->data = ndata;
4
       n->next = node->next; node->next = n;
5
       return 0; // successfully appended
6
   }
7
8
   Node *lx = ... // a linked list
9
   Node *ly = ... // a linked list
10
   for (Node *node = 1x; node != NULL; node = node->next) {
11
       int *dptr = malloc(sizeof(int));
12
       if (!dptr) return;
13
       *dptr = *(node->data);
14
   (-) append_data(ly, dptr); // potential memory-leak
15
   (+) append_data(ly, dptr); free(dptr);
16
    }
17
                                               UAF or DF introduced
```

## cf) SAVER vs. Generative Al

• LLMs do not guarantee safety, e.g., GPT4-generated patch:

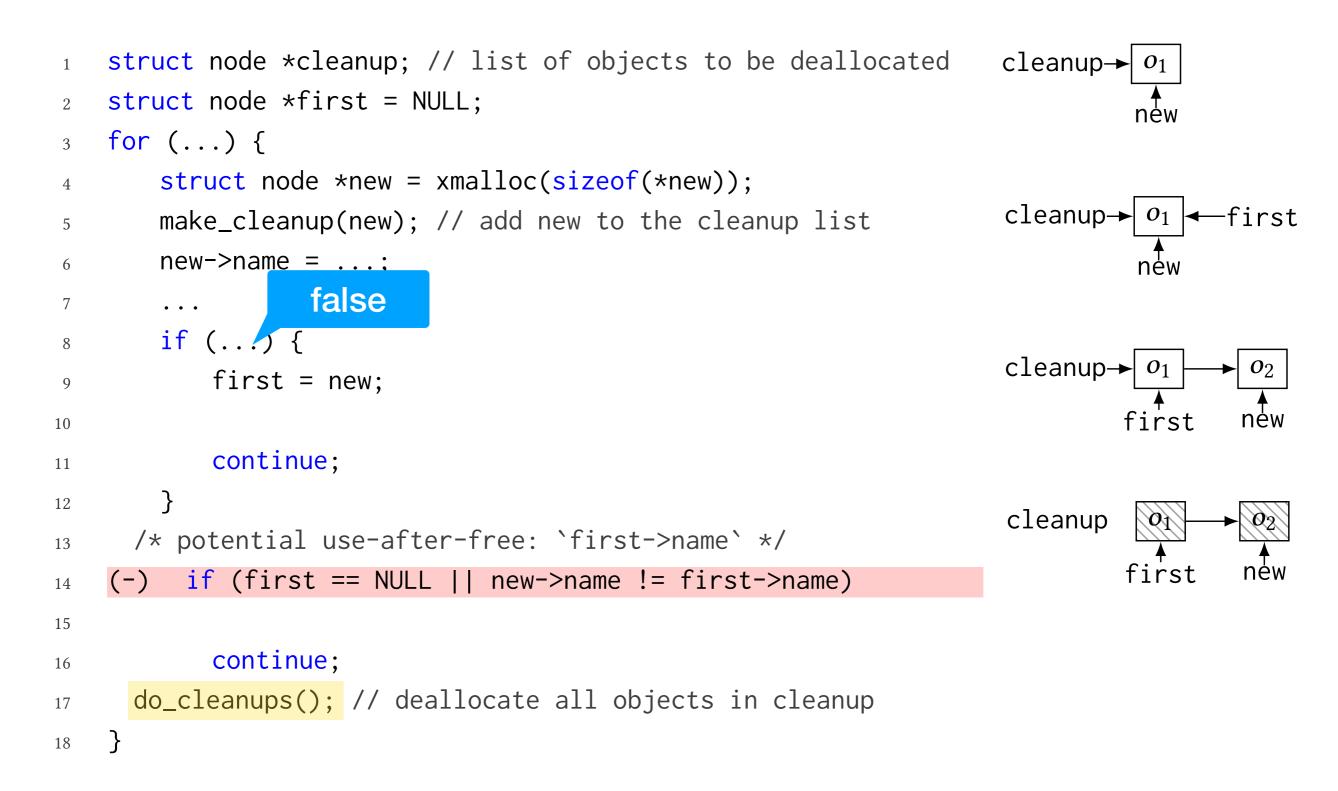
```
int append_data (Node *node, int *ndata) {
1
       if (!(Node *n = malloc(sizeof(Node)))
2
           return -1; // failed to be appended
3
       n->data = ndata;
4
       n->next = node->next; node->next = n;
5
       return 0; // successfully appended
6
   }
7
8
   Node *lx = ... // a linked list
9
   Node *ly = ... // a linked list
10
   for (Node *node = 1x; node != NULL; node = node->next) {
11
       int *dptr = malloc(sizeof(int));
12
       if (!dptr) return;
13
       *dptr = *(node->data);
14
   (-) append_data(ly, dptr); // potential memory-leak
15
   (+) append_data(ly, dptr); free(dptr);
16
    }
17
                                               UAF or DF introduced
```

```
struct node *cleanup; // list of objects to be deallocated
1
   struct node *first = NULL;
2
   for (...) {
3
       struct node *new = xmalloc(sizeof(*new));
4
       make_cleanup(new); // add new to the cleanup list
5
       new->name = ...;
6
7
       . . .
       if (...) {
8
           first = new;
9
10
           continue;
11
       }
12
     /* potential use-after-free: `first->name` */
13
   (-) if (first == NULL || new->name != first->name)
14
15
           continue;
16
     do_cleanups(); // deallocate all objects in cleanup
17
    }
18
```

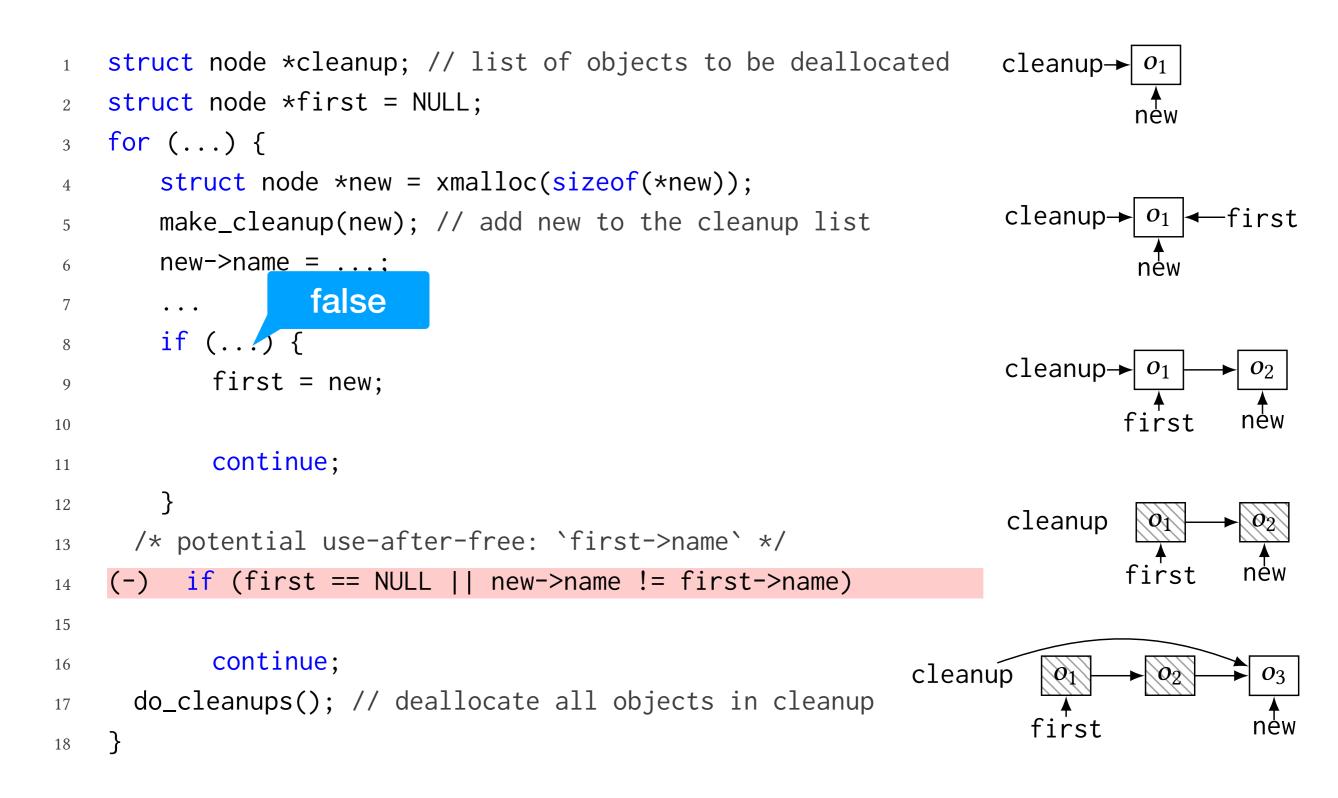
```
struct node *cleanup; // list of objects to be deallocated
                                                                       cleanup→
                                                                                  o<sub>1</sub>
1
    struct node *first = NULL;
2
   for (...) {
3
       struct node *new = xmalloc(sizeof(*new));
4
       make_cleanup(new); // add new to the cleanup list
5
       new->name = ...;
6
7
        . . .
       if (...) {
8
           first = new;
9
10
           continue;
11
        }
12
     /* potential use-after-free: `first->name` */
13
         if (first == NULL || new->name != first->name)
    (-)
14
15
           continue;
16
     do_cleanups(); // deallocate all objects in cleanup
17
    }
18
```

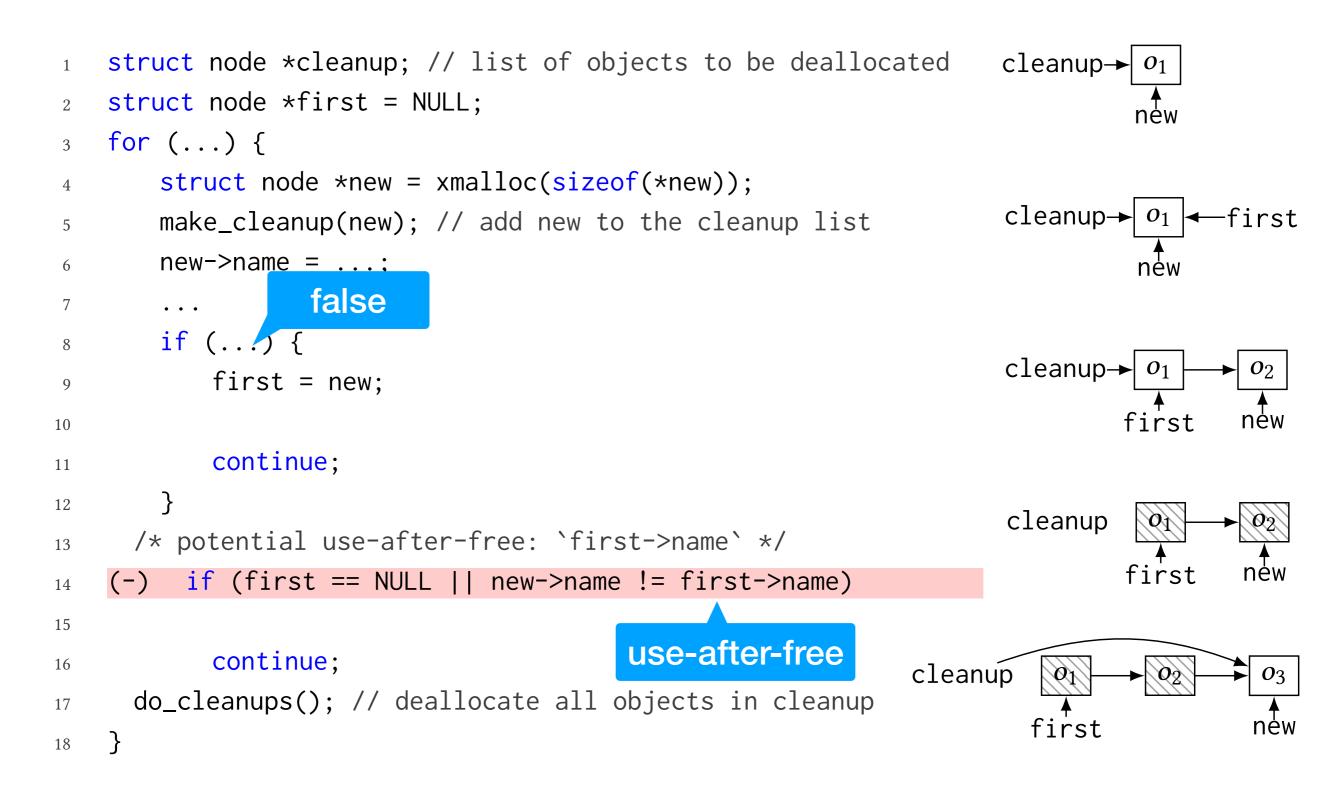
```
struct node *cleanup; // list of objects to be deallocated
                                                                                     o<sub>1</sub>
                                                                           cleanup→
1
    struct node *first = NULL;
2
    for (...) {
3
        struct node *new = xmalloc(sizeof(*new));
4
                                                                           cleanup \rightarrow | o_1 | \leftarrow
                                                                                             first
        make_cleanup(new); // add new to the cleanup list
5
        new->name = ...;
6
                  true
        . . .
7
        if (.../ {
8
            first = new;
9
10
            continue;
11
        }
12
      /* potential use-after-free: `first->name` */
13
          if (first == NULL || new->name != first->name)
    (-)
14
15
            continue;
16
      do_cleanups(); // deallocate all objects in cleanup
17
    }
18
```

```
struct node *cleanup; // list of objects to be deallocated
                                                                            cleanup→
                                                                                       o<sub>1</sub>
1
    struct node *first = NULL;
2
    for (...) {
3
        struct node *new = xmalloc(sizeof(*new));
4
                                                                             cleanup \rightarrow | o_1 | \leftarrow
                                                                                              -first
        make_cleanup(new); // add new to the cleanup list
5
        new->name = ...;
6
7
        . . .
        if (...) {
8
                                                                             cleanup\rightarrow | o_1
            first = new;
9
                                                                                      first
10
            continue;
11
        }
12
      /* potential use-after-free: `first->name` */
13
          if (first == NULL || new->name != first->name)
    (-)
14
15
            continue;
16
      do_cleanups(); // deallocate all objects in cleanup
17
    }
18
```



```
struct node *cleanup; // list of objects to be deallocated
                                                                             cleanup→
                                                                                        o<sub>1</sub>
1
    struct node *first = NULL;
2
    for (...) {
3
        struct node *new = xmalloc(sizeof(*new));
4
                                                                             cleanup \rightarrow | o_1 | \leftarrow
                                                                                               -first
        make_cleanup(new); // add new to the cleanup list
5
        new->name = ...;
6
7
        . . .
        if (...) {
8
                                                                             cleanup\rightarrow | o_1
            first = new;
9
                                                                                       first
                                                                                                 néw
10
            continue;
11
        }
12
                                                                             cleanup
      /* potential use-after-free: `first->name` */
13
                                                                                       first
                                                                                                 nèw
          if (first == NULL || new->name != first->name)
    (-)
14
15
            continue;
16
                                                                     cleanup
                                                                                         Q2
                                                                                                  03
                                                                                 o_1
      do_cleanups(); // deallocate all objects in cleanup
17
                                                                               first
                                                                                                 néw
    }
18
```





## **SAVER-Generated Patch**

```
struct node *cleanup; // list of objects to be deallocated
1
   struct node *first = NULL;
2
   for (...) {
3
       struct node *new = xmalloc(sizeof(*new));
4
       make_cleanup(new); // add new to the cleanup list
5
       new->name = ...;
6
7
       . . .
       if (...) {
8
           first = new;
9
   (+) tmp = first->name;
10
           continue;
11
       }
12
     /* potential use-after-free: `first->name` */
13
   (-) if (first == NULL || new->name != first->name)
14
         if (first == NULL || new->name != tmp)
   (+)
15
           continue;
16
     do_cleanups(); // deallocate all objects in cleanup
17
   }
18
```

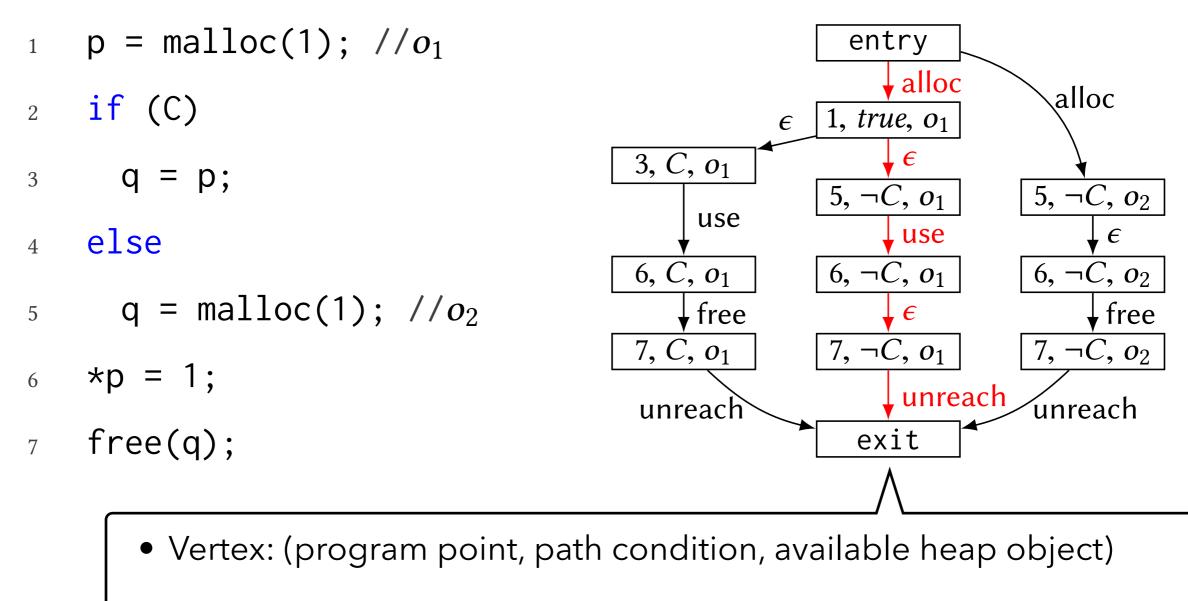
- p = malloc(1);  $//o_1$
- 2 if (C)
- 3 q = p;
- 4 else
- $q = malloc(1); //o_2$
- <sup>6</sup> \*p = 1;
- 7 free(q);

#### Memory leak: $o_1$ is not freed when the false branch is taken

- p = malloc(1);  $//o_1$
- 2 if (C)
- 3 q = p;
- 4 else
- $q = malloc(1); //o_2$
- 6 \*p = 1; 7 free(q);

#### Memory leak: $o_1$ is not freed when the false branch is taken

1. Run a static analysis to generate object flow graph

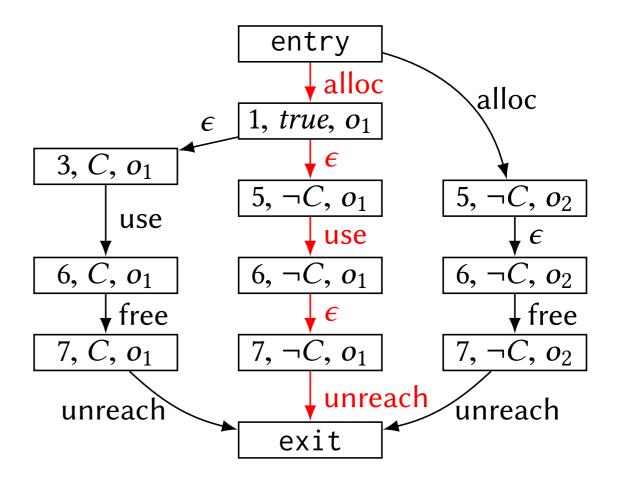


• Edge: control flow labeled w/ events that could occur for objects

2. Relabel object flow graph to eliminate buggy paths

- 1 p = malloc(1); //o<sub>1</sub>
  2 if (C)
  3 q = p;
- 4 else
- $q = malloc(1); //o_2$
- $^{6}$  \*p = 1;

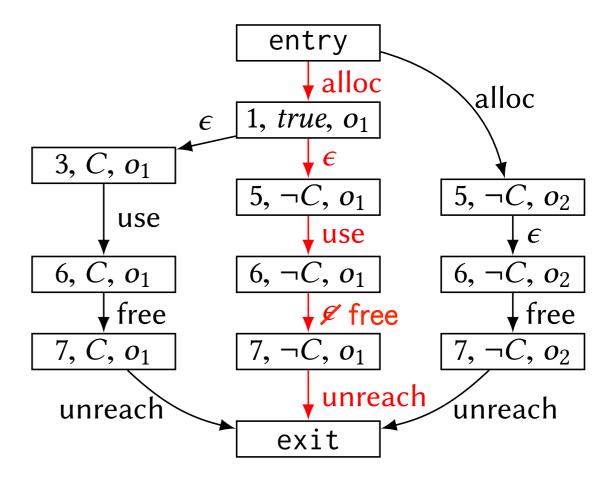
7 free(q);



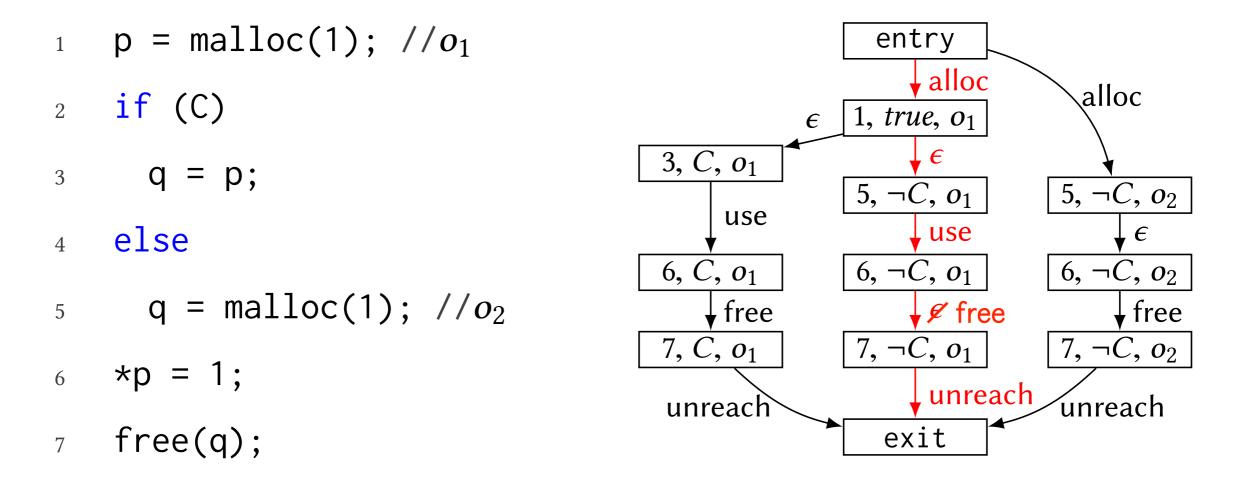
2. Relabel object flow graph to eliminate buggy paths

- 1 p = malloc(1); //o<sub>1</sub>
  2 if (C)
  3 q = p;
- 4 else
- $q = malloc(1); //o_2$
- <sup>6</sup> \*p = 1;

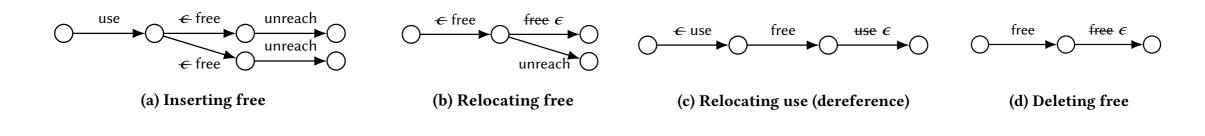
7 free(q);



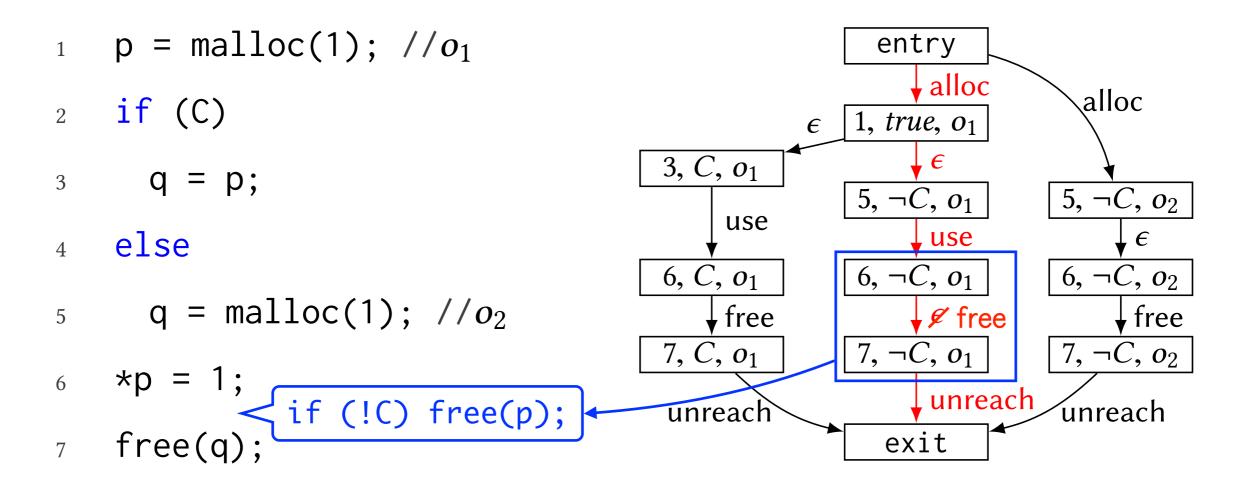
2. Relabel object flow graph to eliminate buggy paths



• SAVER supports four types of re-labeling strategies:



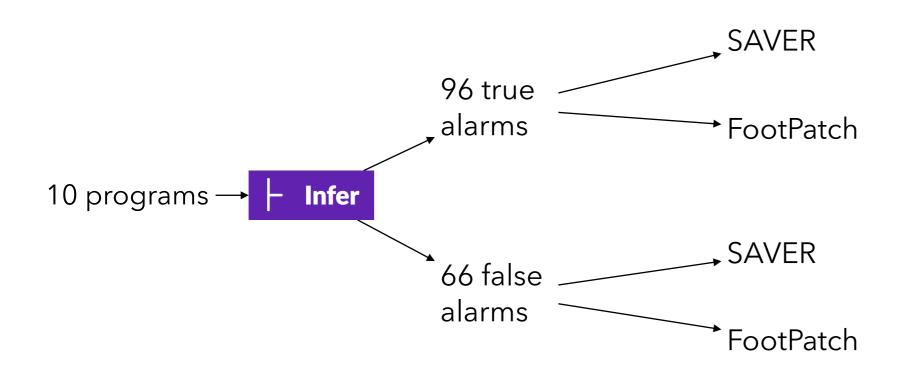
3. Generate a patch from the re-labeled edge



# For Deployment

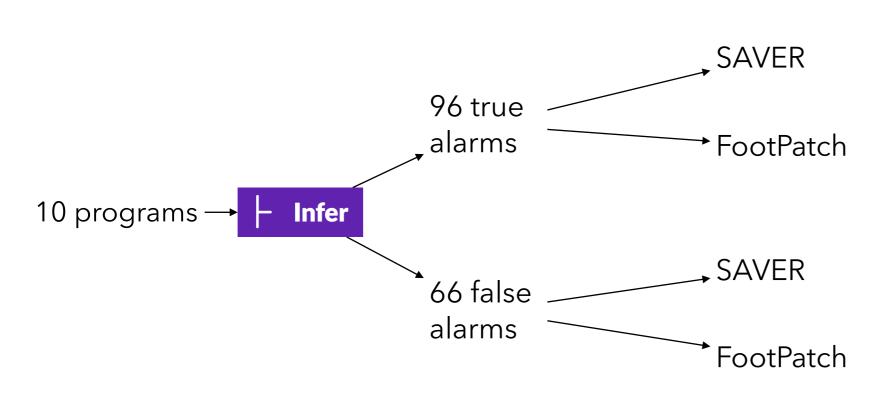
- Improved scalability (in the paper)
  - Selective path sensitivity
  - Program slicing
- Improved usability (not in the paper)
  - Build failures
  - Robust translation from IR-level patches to source-level

		Ini	FER			S	AVER						F	оотРа	тсн [	60]		
Program	kLoC	#T	#F	Pre(s)	Fix(s)	GT	✓ <sub>T</sub>	$\Delta_{\mathbf{T}}$	X <sub>T</sub>	G <sub>F</sub>	X <sub>F</sub>	Fix(s)	G <sub>T</sub>	✓ <sub>T</sub>	$\Delta_{\mathbf{T}}$	X <sub>T</sub>	G <sub>F</sub>	× <sub>F</sub>
rappel (ad8efd7)	2.2	1	0	2.2	0.0	1	1	0	0	0	0	8.9	1	1	0	0	0	0
flex (d3de49f)	22.3	3	4	26.3	2.5	0	0	0	0	0	0	51.0	0	0	0	0	1	1
WavPack (22977b2)	31.2	1	2	44.6	22.1	0	0	0	0	0	0	67.9	0	0	0	0	2	2
Swoole (a4256e4)	43.0	15	3	88.5	10.1	11	11	0	0	0	0	392.5	9	7	0	2	1	1
lxc (72cc48f)	49.9	3	5	230.6	5.8	3	3	0	0	0	0	179.6	0	0	0	0	1	1
p11-kit (ead7ara)	62.9	33	9	646.2	288.8	24	24	0	0	0	0	566.4	8	7	1	0	2	2
x264 (d4099dd)	73.2	10	0	144.3	9.9	10	10	0	0	0	0	426.9	2	2	0	0	0	0
recutils-1.8	92.0	10	11	144.1	44.4	8	8	0	0	0	0	662.2	3	2	1	0	0	0
inetutils-1.9.4	116.9	4	5	44.8	2.5	4	4	0	0	0	0	182.1	0	0	0	0	0	0
snort-2.9.13	320.8	16	27	2372.0	216.0	11	10	1	0	0	0	4636.4	3	0	0	3	19	18
Total	814.4	96	66	3743.6	602.1	72	71	1	0	0	0	7173.9	26	19	2	5	26	25



Existing memory error repair tool [ICSE 2018]

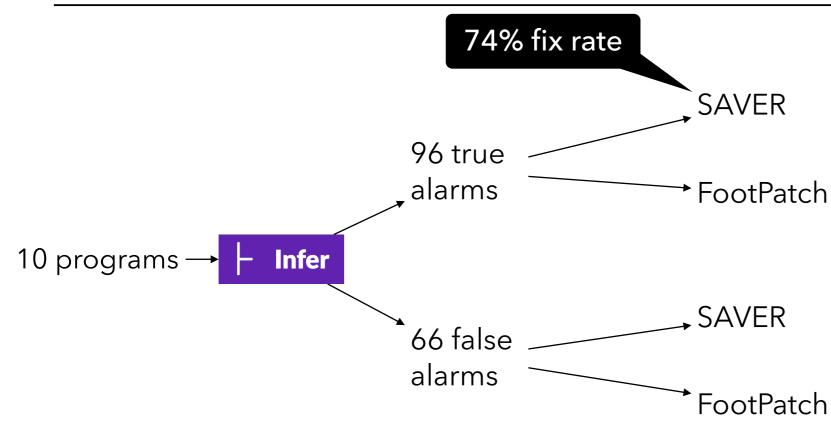
		Inf	FER			S	AVER						F	οοτ <b>γ</b>	тсн [	[60]		
Program	kLoC	#T	#F	Pre(s)	Fix(s)	GT	✓T	$\Delta_{\mathbf{T}}$	X <sub>T</sub>	G <sub>F</sub>	X <sub>F</sub>	Fix(s)	G <sub>T</sub>	✓T	$\Delta_{\mathbf{T}}$	X <sub>T</sub>	G <sub>F</sub>	<b>X</b> <sub>F</sub>
rappel (ad8efd7)	2.2	1	0	2.2	0.0	1	1	0	0	0	0	8.9	1	1	0	0	0	0
flex (d3de49f)	22.3	3	4	26.3	2.5	0	0	0	0	0	0	51.0	0	0	0	0	1	1
WavPack (22977b2)	31.2	1	2	44.6	22.1	0	0	0	0	0	0	67.9	0	0	0	0	2	2
Swoole (a4256e4)	43.0	15	3	88.5	10.1	11	11	0	0	0	0	392.5	9	7	0	2	1	1
lxc (72cc48f)	49.9	3	5	230.6	5.8	3	3	0	0	0	0	179.6	0	0	0	0	1	1
p11-kit (ead7ara)	62.9	33	9	646.2	288.8	24	24	0	0	0	0	566.4	8	7	1	0	2	2
x264 (d4099dd)	73.2	10	0	144.3	9.9	10	10	0	0	0	0	426.9	2	2	0	0	0	0
recutils-1.8	92.0	10	11	144.1	44.4	8	8	0	0	0	0	662.2	3	2	1	0	0	0
inetutils-1.9.4	116.9	4	5	44.8	2.5	4	4	0	0	0	0	182.1	0	0	0	0	0	0
snort-2.9.13	320.8	16	27	2372.0	216.0	11	10	1	0	0	0	4636.4	3	0	0	3	19	18
Total	814.4	96	66	3743.6	602.1	72	71	1	0	0	0	7173.9	26	19	2	5	26	25



Generated 72 patches (Fixed: 71, Safe: 1, Unsafe: 0)

Existing memory error repair tool [ICSE 2018]

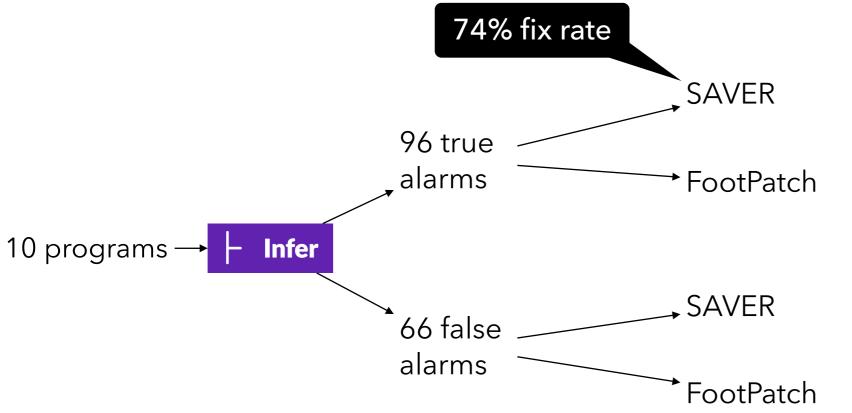
		Ini	FER			S	AVER	L					F	οοτ <b>γ</b>	атсн [	[60]		
Program	kLoC	#T	#F	Pre(s)	Fix(s)	GT	✓ <sub>T</sub>	$\Delta_{\mathbf{T}}$	X <sub>T</sub>	G <sub>F</sub>	$\boldsymbol{X}_{\mathrm{F}}$	Fix(s)	G <sub>T</sub>	✓ <sub>T</sub>	$\Delta_{\mathbf{T}}$	X <sub>T</sub>	G <sub>F</sub>	$\boldsymbol{X}_{\mathrm{F}}$
rappel (ad8efd7)	2.2	1	0	2.2	0.0	1	1	0	0	0	0	8.9	1	1	0	0	0	0
flex (d3de49f)	22.3	3	4	26.3	2.5	0	0	0	0	0	0	51.0	0	0	0	0	1	1
WavPack (22977b2)	31.2	1	2	44.6	22.1	0	0	0	0	0	0	67.9	0	0	0	0	2	2
Swoole (a4256e4)	43.0	15	3	88.5	10.1	11	11	0	0	0	0	392.5	9	7	0	2	1	1
lxc (72cc48f)	49.9	3	5	230.6	5.8	3	3	0	0	0	0	179.6	0	0	0	0	1	1
p11-kit (ead7ara)	62.9	33	9	646.2	288.8	24	24	0	0	0	0	566.4	8	7	1	0	2	2
x264 (d4099dd)	73.2	10	0	144.3	9.9	10	10	0	0	0	0	426.9	2	2	0	0	0	0
recutils-1.8	92.0	10	11	144.1	44.4	8	8	0	0	0	0	662.2	3	2	1	0	0	0
inetutils-1.9.4	116.9	4	5	44.8	2.5	4	4	0	0	0	0	182.1	0	0	0	0	0	0
snort-2.9.13	320.8	16	27	2372.0	216.0	11	10	1	0	0	0	4636.4	3	0	0	3	19	18
Total	814.4	96	66	3743.6	602.1	72	71	1	0	0	0	7173.9	26	19	2	5	26	25



Generated 72 patches (Fixed: 71, Safe: 1, Unsafe: 0)

Existing memory error repair tool [ICSE 2018]

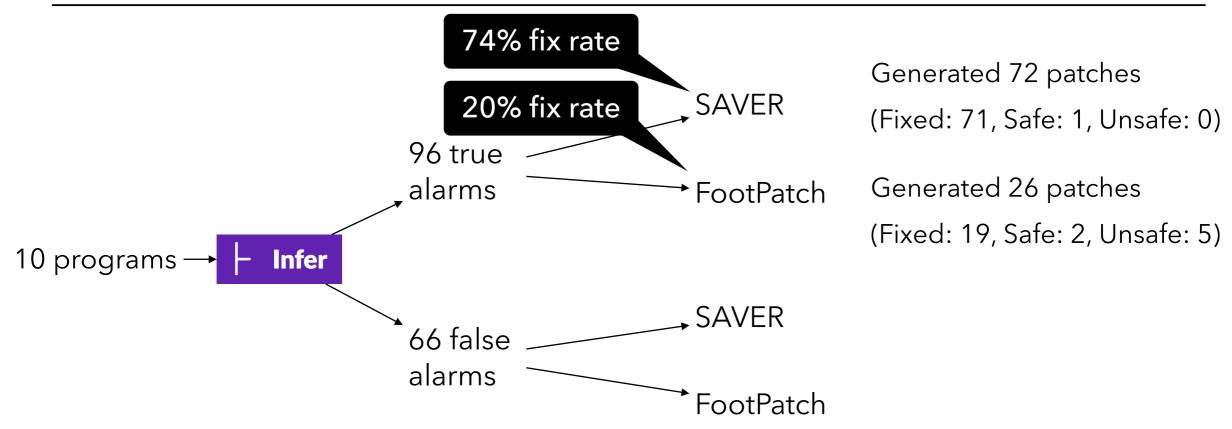
		Ini	FER			S	AVER	R					F	оотРа	тсн [	[60]		
Program	kLoC	#T	#F	Pre(s)	Fix(s)	GT	✓T	$\Delta_{\mathbf{T}}$	X <sub>T</sub>	G <sub>F</sub>	X <sub>F</sub>	Fix(s)	G <sub>T</sub>	✓ <sub>T</sub>	$\Delta_{\mathbf{T}}$	X <sub>T</sub>	G <sub>F</sub>	× <sub>F</sub>
rappel (ad8efd7)	2.2	1	0	2.2	0.0	1	1	0	0	0	0	8.9	1	1	0	0	0	0
flex (d3de49f)	22.3	3	4	26.3	2.5	0	0	0	0	0	0	51.0	0	0	0	0	1	1
WavPack (22977b2)	31.2	1	2	44.6	22.1	0	0	0	0	0	0	67.9	0	0	0	0	2	2
Swoole (a4256e4)	43.0	15	3	88.5	10.1	11	11	0	0	0	0	392.5	9	7	0	2	1	1
lxc (72cc48f)	49.9	3	5	230.6	5.8	3	3	0	0	0	0	179.6	0	0	0	0	1	1
p11-kit (ead7ara)	62.9	33	9	646.2	288.8	24	24	0	0	0	0	566.4	8	7	1	0	2	2
x264 (d4099dd)	73.2	10	0	144.3	9.9	10	10	0	0	0	0	426.9	2	2	0	0	0	0
recutils-1.8	92.0	10	11	144.1	44.4	8	8	0	0	0	0	662.2	3	2	1	0	0	0
inetutils-1.9.4	116.9	4	5	44.8	2.5	4	4	0	0	0	0	182.1	0	0	0	0	0	0
snort-2.9.13	320.8	16	27	2372.0	216.0	11	10	1	0	0	0	4636.4	3	0	0	3	19	18
Total	814.4	96	66	3743.6	602.1	72	71	1	0	0	0	7173.9	26	19	2	5	26	25



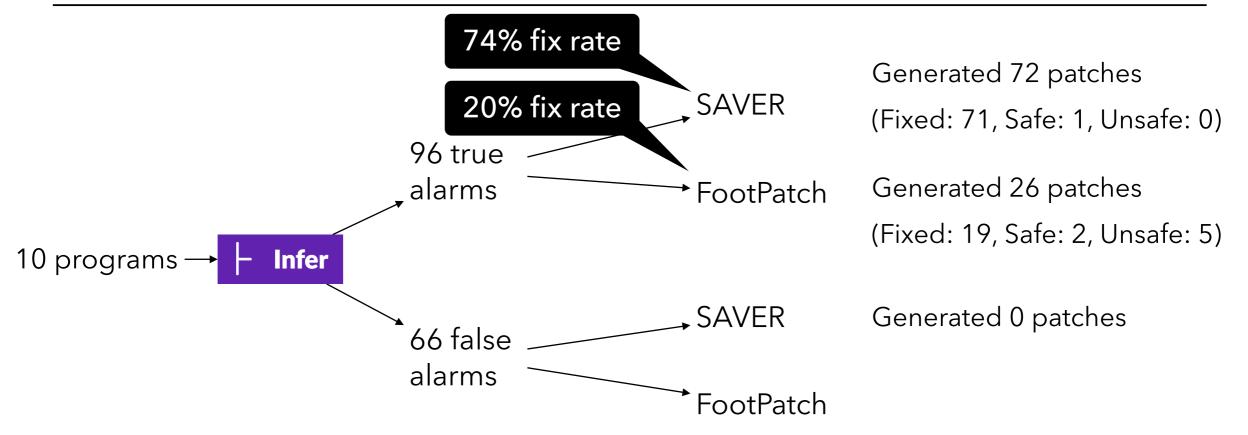
Generated 72 patches (Fixed: 71, Safe: 1, Unsafe: 0) Generated 26 patches

(Fixed: 19, Safe: 2, Unsafe: 5)

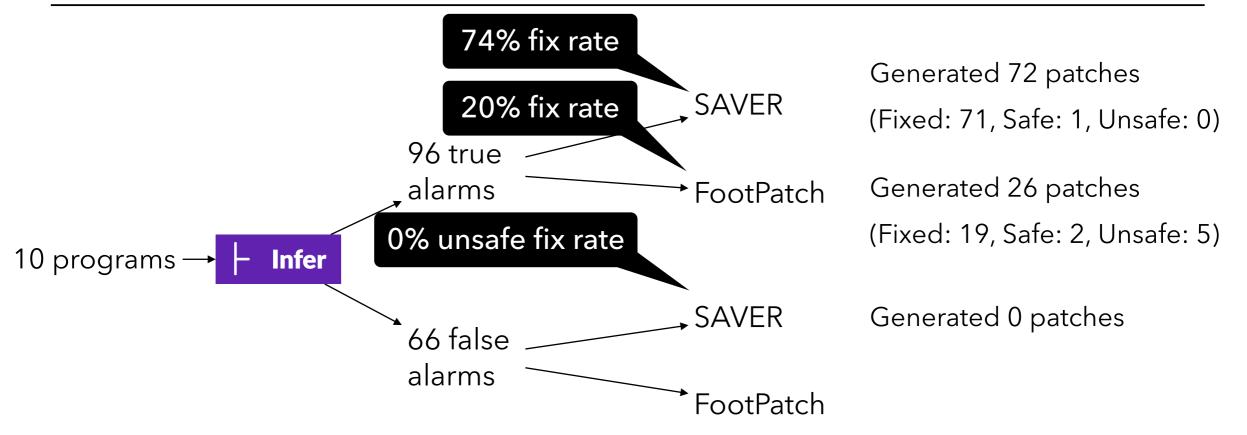
		Ini	FER			S	AVER	ł					F	оотРа	тсн [	[60]		
Program	kLoC	#T	#F	Pre(s)	Fix(s)	GT	✓ <sub>T</sub>	$\Delta_{\mathbf{T}}$	X <sub>T</sub>	G <sub>F</sub>	X <sub>F</sub>	Fix(s)	G <sub>T</sub>	✓ <sub>T</sub>	$\Delta_{\mathbf{T}}$	X <sub>T</sub>	G <sub>F</sub>	$\boldsymbol{X}_{\mathrm{F}}$
rappel (ad8efd7)	2.2	1	0	2.2	0.0	1	1	0	0	0	0	8.9	1	1	0	0	0	0
flex (d3de49f)	22.3	3	4	26.3	2.5	0	0	0	0	0	0	51.0	0	0	0	0	1	1
WavPack (22977b2)	31.2	1	2	44.6	22.1	0	0	0	0	0	0	67.9	0	0	0	0	2	2
Swoole (a4256e4)	43.0	15	3	88.5	10.1	11	11	0	0	0	0	392.5	9	7	0	2	1	1
lxc (72cc48f)	49.9	3	5	230.6	5.8	3	3	0	0	0	0	179.6	0	0	0	0	1	1
p11-kit (ead7ara)	62.9	33	9	646.2	288.8	24	24	0	0	0	0	566.4	8	7	1	0	2	2
x264 (d4099dd)	73.2	10	0	144.3	9.9	10	10	0	0	0	0	426.9	2	2	0	0	0	0
recutils-1.8	92.0	10	11	144.1	44.4	8	8	0	0	0	0	662.2	3	2	1	0	0	0
inetutils-1.9.4	116.9	4	5	44.8	2.5	4	4	0	0	0	0	182.1	0	0	0	0	0	0
snort-2.9.13	320.8	16	27	2372.0	216.0	11	10	1	0	0	0	4636.4	3	0	0	3	19	18
Total	814.4	96	66	3743.6	602.1	72	71	1	0	0	0	7173.9	26	19	2	5	26	25



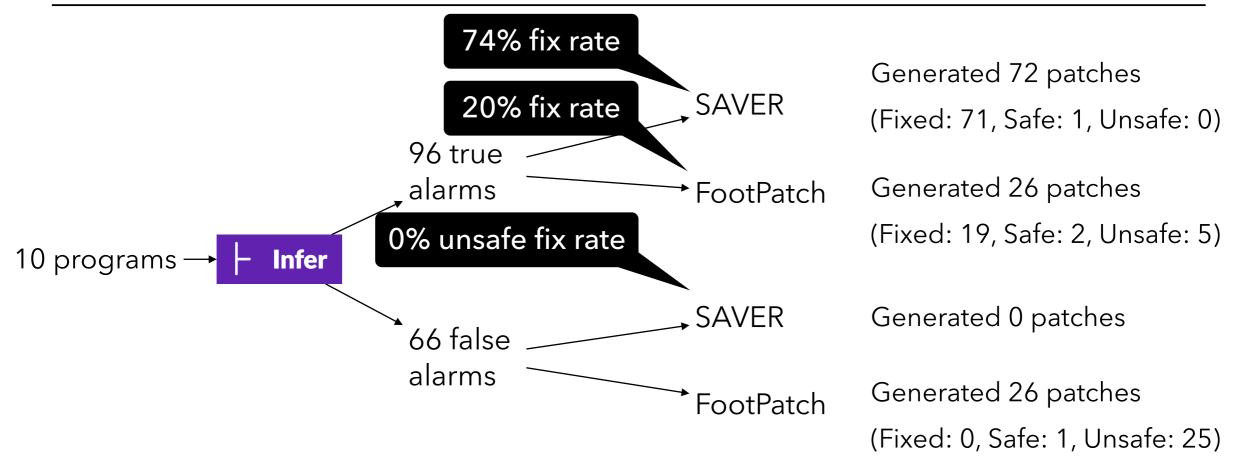
		Ini	FER			S	AVER	R					F	оотРа	тсн [	[60]		
Program	kLoC	#T	#F	Pre(s)	Fix(s)	GT	✓T	$\Delta_{\mathbf{T}}$	X <sub>T</sub>	G <sub>F</sub>	$\boldsymbol{X}_{\mathrm{F}}$	Fix(s)	G <sub>T</sub>	✓ <sub>T</sub>	$\Delta_{\mathbf{T}}$	X <sub>T</sub>	G <sub>F</sub>	× <sub>F</sub>
rappel (ad8efd7)	2.2	1	0	2.2	0.0	1	1	0	0	0	0	8.9	1	1	0	0	0	0
flex (d3de49f)	22.3	3	4	26.3	2.5	0	0	0	0	0	0	51.0	0	0	0	0	1	1
WavPack (22977b2)	31.2	1	2	44.6	22.1	0	0	0	0	0	0	67.9	0	0	0	0	2	2
Swoole (a4256e4)	43.0	15	3	88.5	10.1	11	11	0	0	0	0	392.5	9	7	0	2	1	1
lxc (72cc48f)	49.9	3	5	230.6	5.8	3	3	0	0	0	0	179.6	0	0	0	0	1	1
p11-kit (ead7ara)	62.9	33	9	646.2	288.8	24	24	0	0	0	0	566.4	8	7	1	0	2	2
x264 (d4099dd)	73.2	10	0	144.3	9.9	10	10	0	0	0	0	426.9	2	2	0	0	0	0
recutils-1.8	92.0	10	11	144.1	44.4	8	8	0	0	0	0	662.2	3	2	1	0	0	0
inetutils-1.9.4	116.9	4	5	44.8	2.5	4	4	0	0	0	0	182.1	0	0	0	0	0	0
snort-2.9.13	320.8	16	27	2372.0	216.0	11	10	1	0	0	0	4636.4	3	0	0	3	19	18
Total	814.4	96	66	3743.6	602.1	72	71	1	0	0	0	7173.9	26	19	2	5	26	25



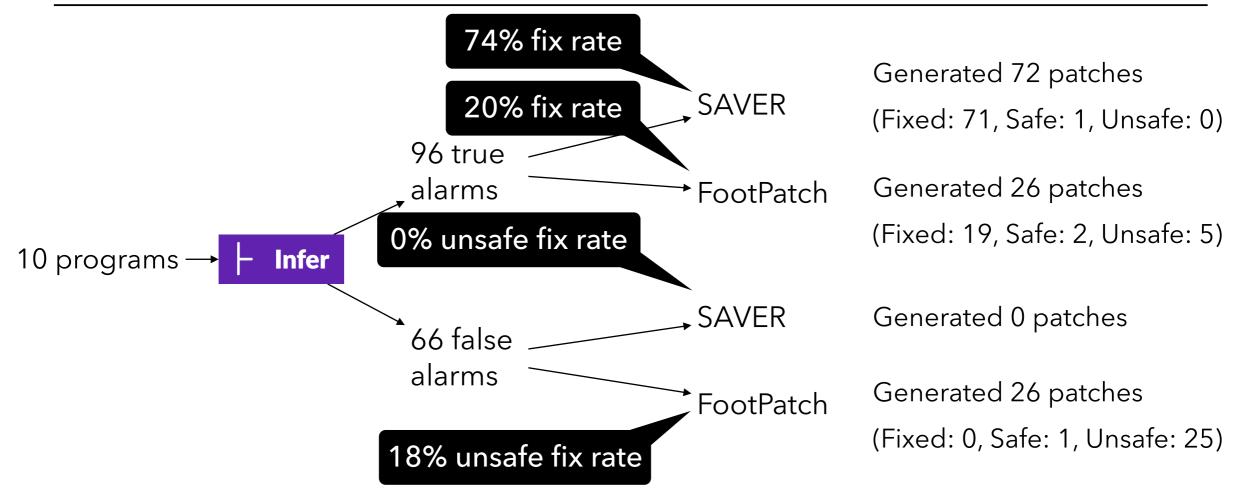
		Ini	FER			S	AVER	2					F	оотРа	тсн [	60]		
Program	kLoC	#T	#F	Pre(s)	Fix(s)	GT	✓ <sub>T</sub>	$\Delta_{\mathbf{T}}$	X <sub>T</sub>	G <sub>F</sub>	X <sub>F</sub>	Fix(s)	G <sub>T</sub>	✓ <sub>T</sub>	$\Delta_{\mathbf{T}}$	X <sub>T</sub>	G <sub>F</sub>	× <sub>F</sub>
rappel (ad8efd7)	2.2	1	0	2.2	0.0	1	1	0	0	0	0	8.9	1	1	0	0	0	0
flex (d3de49f)	22.3	3	4	26.3	2.5	0	0	0	0	0	0	51.0	0	0	0	0	1	1
WavPack (22977b2)	31.2	1	2	44.6	22.1	0	0	0	0	0	0	67.9	0	0	0	0	2	2
Swoole (a4256e4)	43.0	15	3	88.5	10.1	11	11	0	0	0	0	392.5	9	7	0	2	1	1
lxc (72cc48f)	49.9	3	5	230.6	5.8	3	3	0	0	0	0	179.6	0	0	0	0	1	1
p11-kit (ead7ara)	62.9	33	9	646.2	288.8	24	24	0	0	0	0	566.4	8	7	1	0	2	2
x264 (d4099dd)	73.2	10	0	144.3	9.9	10	10	0	0	0	0	426.9	2	2	0	0	0	0
recutils-1.8	92.0	10	11	144.1	44.4	8	8	0	0	0	0	662.2	3	2	1	0	0	0
inetutils-1.9.4	116.9	4	5	44.8	2.5	4	4	0	0	0	0	182.1	0	0	0	0	0	0
snort-2.9.13	320.8	16	27	2372.0	216.0	11	10	1	0	0	0	4636.4	3	0	0	3	19	18
Total	814.4	96	66	3743.6	602.1	72	71	1	0	0	0	7173.9	26	19	2	5	26	25



		Ini	FER			S	AVER	Ł					F	оотРа	тсн [	[60]		
Program	kLoC	#T	#F	Pre(s)	Fix(s)	GT	✓T	$\Delta_{\mathbf{T}}$	X <sub>T</sub>	G <sub>F</sub>	× <sub>F</sub>	Fix(s)	G <sub>T</sub>	✓ <sub>T</sub>	$\Delta_{\mathbf{T}}$	X <sub>T</sub>	G <sub>F</sub>	× <sub>F</sub>
rappel (ad8efd7)	2.2	1	0	2.2	0.0	1	1	0	0	0	0	8.9	1	1	0	0	0	0
flex (d3de49f)	22.3	3	4	26.3	2.5	0	0	0	0	0	0	51.0	0	0	0	0	1	1
WavPack (22977b2)	31.2	1	2	44.6	22.1	0	0	0	0	0	0	67.9	0	0	0	0	2	2
Swoole (a4256e4)	43.0	15	3	88.5	10.1	11	11	0	0	0	0	392.5	9	7	0	2	1	1
lxc (72cc48f)	49.9	3	5	230.6	5.8	3	3	0	0	0	0	179.6	0	0	0	0	1	1
p11-kit (ead7ara)	62.9	33	9	646.2	288.8	24	24	0	0	0	0	566.4	8	7	1	0	2	2
x264 (d4099dd)	73.2	10	0	144.3	9.9	10	10	0	0	0	0	426.9	2	2	0	0	0	0
recutils-1.8	92.0	10	11	144.1	44.4	8	8	0	0	0	0	662.2	3	2	1	0	0	0
inetutils-1.9.4	116.9	4	5	44.8	2.5	4	4	0	0	0	0	182.1	0	0	0	0	0	0
snort-2.9.13	320.8	16	27	2372.0	216.0	11	10	1	0	0	0	4636.4	3	0	0	3	19	18
Total	814.4	96	66	3743.6	602.1	72	71	1	0	0	0	7173.9	26	19	2	5	26	25



		Ini	FER			S	AVER	Ł					F	оотРа	атсн [	60]		
Program	kLoC	#T	#F	Pre(s)	Fix(s)	GT	✓ <sub>T</sub>	$\Delta_{\mathbf{T}}$	X <sub>T</sub>	G <sub>F</sub>	$\boldsymbol{X}_{\mathrm{F}}$	Fix(s)	G <sub>T</sub>	✓ <sub>T</sub>	$\Delta_{\mathbf{T}}$	X <sub>T</sub>	G <sub>F</sub>	$\boldsymbol{X}_{\mathrm{F}}$
rappel (ad8efd7)	2.2	1	0	2.2	0.0	1	1	0	0	0	0	8.9	1	1	0	0	0	0
flex (d3de49f)	22.3	3	4	26.3	2.5	0	0	0	0	0	0	51.0	0	0	0	0	1	1
WavPack (22977b2)	31.2	1	2	44.6	22.1	0	0	0	0	0	0	67.9	0	0	0	0	2	2
Swoole (a4256e4)	43.0	15	3	88.5	10.1	11	11	0	0	0	0	392.5	9	7	0	2	1	1
lxc (72cc48f)	49.9	3	5	230.6	5.8	3	3	0	0	0	0	179.6	0	0	0	0	1	1
p11-kit (ead7ara)	62.9	33	9	646.2	288.8	24	24	0	0	0	0	566.4	8	7	1	0	2	2
x264 (d4099dd)	73.2	10	0	144.3	9.9	10	10	0	0	0	0	426.9	2	2	0	0	0	0
recutils-1.8	92.0	10	11	144.1	44.4	8	8	0	0	0	0	662.2	3	2	1	0	0	0
inetutils-1.9.4	116.9	4	5	44.8	2.5	4	4	0	0	0	0	182.1	0	0	0	0	0	0
snort-2.9.13	320.8	16	27	2372.0	216.0	11	10	1	0	0	0	4636.4	3	0	0	3	19	18
Total	814.4	96	66	3743.6	602.1	72	71	1	0	0	0	7173.9	26	19	2	5	26	25



# Summary

- Static analysis-based program repair for C memory errors
  - Scalable, precise, and safe
  - Successfully deployed in industry
  - <u>https://github.com/kupl/kaprese</u>

Thank you!