

# CVO103: Programming Languages

## Lecture 2 — Inductive Definitions (2)

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# Contents

- More examples of inductive definitions
  - ▶ natural numbers, strings, booleans
  - ▶ lists, binary trees
  - ▶ arithmetic expressions, propositional logic
- Structural induction
  - ▶ three example proofs

# Natural Numbers

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$$\mathbb{N} = \{0, 1, 2, 3, \dots\}$$

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is inductively defined:

$$\bar{0} \quad \frac{n}{n+1}$$

The inference rules can be expressed by a grammar:

$$n \rightarrow 0 \mid n + 1$$

Interpretation:

- 0 is a natural number.
- If  $n$  is a natural number then so is  $n + 1$ .

# Strings

The set of strings over alphabet  $\{a, \dots, z\}$ , e.g.,  $\epsilon$ ,  $a$ ,  $b$ ,  $\dots$ ,  $z$ ,  $aa$ ,  $ab$ ,  $\dots$ ,  $az$ ,  $ba$ ,  $\dots$ ,  $az$ ,  $aaa$ ,  $\dots$ ,  $zzz$ , and so on.

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or simply,

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In grammar:

$$\begin{array}{l} \alpha \rightarrow \epsilon \\ \quad | \quad x\alpha \quad (x \in \{a, \dots, z\}) \end{array}$$

# Boolean Values

The set of boolean values:

$$\mathbb{B} = \{\mathit{true}, \mathit{false}\}.$$

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$$\mathbb{B} = \{true, false\}.$$

If a set is finite, just enumerate all of its elements by axioms:

$$\overline{true} \quad \overline{false}$$

In grammar:

$$b \rightarrow true \mid false$$

# Lists

Examples of lists of integers:

- 1 **nil**
- 2 **14 · nil**
- 3 **3 · 14 · nil**
- 4 **-7 · 3 · 14 · nil**

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Inference rules:

$$\frac{}{\mathbf{nil}} \quad \frac{l}{n \cdot l} \quad n \in \mathbb{Z}$$

In grammar:

$$l \rightarrow \begin{array}{l} \mathbf{nil} \\ | \\ n \cdot l \quad (n \in \mathbb{Z}) \end{array}$$

## Lists

A proof that  $-7 \cdot 3 \cdot 14 \cdot \mathbf{nil}$  is a list of integers:

$$\frac{\frac{\frac{\mathbf{nil}}{14 \cdot \mathbf{nil}}}{3 \cdot 14 \cdot \mathbf{nil}}}{-7 \cdot 3 \cdot 14 \cdot \mathbf{nil}} \quad \begin{array}{l} 14 \in \mathbb{Z} \\ 3 \in \mathbb{Z} \\ -7 \in \mathbb{Z} \end{array}$$

The proof tree is also called *derivation tree* or *deduction tree*.

# Binary Trees

Examples of binary trees:

- ① leaf
- ② (2, leaf, leaf)
- ③ (1, (2, leaf, leaf), leaf)
- ④ (1, (2, leaf, leaf), (3, (4, leaf, leaf), leaf))

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- 1 **leaf**
- 2 **(2, leaf, leaf)**
- 3 **(1, (2, leaf, leaf), leaf)**
- 4 **(1, (2, leaf, leaf), (3, (4, leaf, leaf), leaf))**

Inference rules:

$$\frac{}{\mathbf{leaf}} \quad \frac{t_1 \quad t_2}{(n, t_1, t_2)} \quad n \in \mathbb{Z}$$

In grammar:

$$t \rightarrow \mathbf{leaf} \\ | \quad (n, t, t) \quad (n \in \mathbb{Z})$$



# Binary Trees

A proof that

$(1, (2, \text{leaf}, \text{leaf}), (3, (4, \text{leaf}, \text{leaf}), \text{leaf}))$

is a binary tree:

$$\frac{\frac{\overline{\text{leaf}}}{(2, \text{leaf}, \text{leaf})} \quad 2 \in \mathbb{Z} \quad \frac{\frac{\overline{\text{leaf}}}{(4, \text{leaf}, \text{leaf})} \quad 4 \in \mathbb{Z}}{(3, (4, \text{leaf}, \text{leaf}), \text{leaf})} \quad 3 \in \mathbb{Z}}{(1, (2, \text{leaf}, \text{leaf}), (3, (4, \text{leaf}, \text{leaf}), \text{leaf}))} \quad 1 \in \mathbb{Z}$$

## Binary Trees: a different version

Binary tree examples:  $1$ ,  $(1, \mathbf{nil})$ ,  $(1, 2)$ ,  $((1, 2), \mathbf{nil})$ ,  $((1, 2), (3, 4))$ .

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Inference rules:

$$\overline{n} \quad n \in \mathbb{Z} \qquad \frac{t}{(t, \mathbf{nil})} \qquad \frac{t}{(\mathbf{nil}, t)} \qquad \frac{\overline{t_1} \quad \overline{t_2}}{\overline{(t_1, t_2)}}$$

In grammar:

$$\begin{array}{l} t \rightarrow n \quad (n \in \mathbb{Z}) \\ \quad | \quad (t, \mathbf{nil}) \\ \quad | \quad (\mathbf{nil}, t) \\ \quad | \quad (t, t) \end{array}$$

A proof that  $((1, 2), (3, \mathbf{nil}))$  is a binary tree:

$$\frac{\frac{\overline{1} \quad \overline{2}}{\overline{(1, 2)}} \quad \frac{\overline{3}}{\overline{(3, \mathbf{nil})}}}{\overline{((1, 2), (3, \mathbf{nil}))}}$$

# Expressions

Expression examples:  $2$ ,  $-2$ ,  $1 + 2$ ,  $1 + (2 * (-3))$ , etc.

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Inference rules:

$$\overline{n} \quad n \in \mathbb{Z} \quad \frac{e}{-e} \quad \frac{e_1 \quad e_2}{e_1 + e_2} \quad \frac{e_1 \quad e_2}{e_1 * e_2} \quad \frac{e}{(e)}$$

In grammar:

$$e \rightarrow \begin{array}{l} n \quad (n \in \mathbb{Z}) \\ -e \\ e + e \\ e * e \\ (e) \end{array}$$

Example:

$$\frac{\overline{1} \quad \frac{\overline{2} \quad \frac{\overline{3} \quad \overline{-3}}{\overline{-3}}}{\overline{2} \quad \overline{(-3)}}}{\overline{2 * (-3)}}}{\overline{1} \quad \overline{(2 * (-3))}}}{\overline{1 + (2 * (-3))}}$$

# Propositional Logic

Examples:

- $T, F$
- $T \wedge F$
- $T \vee F$
- $(T \wedge F) \wedge (T \vee F)$
- $T \Rightarrow (F \Rightarrow T)$

# Propositional Logic

Syntax:

$$\begin{array}{l} f \rightarrow T \mid F \\ | \neg f \\ | f \wedge f \\ | f \vee f \\ | f \Rightarrow f \end{array}$$

Semantics ( $\llbracket f \rrbracket$ ):

$$\begin{array}{l} \llbracket T \rrbracket = \text{true} \\ \llbracket F \rrbracket = \text{false} \\ \llbracket \neg f \rrbracket = \text{not } \llbracket f \rrbracket \\ \llbracket f_1 \wedge f_2 \rrbracket = \llbracket f_1 \rrbracket \text{ and also } \llbracket f_2 \rrbracket \\ \llbracket f_1 \vee f_2 \rrbracket = \llbracket f_1 \rrbracket \text{ or else } \llbracket f_2 \rrbracket \\ \llbracket f_1 \Rightarrow f_2 \rrbracket = \llbracket f_1 \rrbracket \text{ implies } \llbracket f_2 \rrbracket \end{array}$$

# Propositional Logic

$$\begin{aligned} \llbracket (T \wedge (T \vee F)) \Rightarrow F \rrbracket &= \llbracket T \wedge (T \vee F) \rrbracket \text{ implies } \llbracket F \rrbracket \\ &= (\llbracket T \rrbracket \text{ andalso } \llbracket T \vee F \rrbracket) \text{ implies } \textit{false} \\ &= (\textit{true} \text{ andalso } (\llbracket T \rrbracket \text{ orelse } \llbracket F \rrbracket)) \text{ implies } \textit{false} \\ &= (\textit{true} \text{ andalso } (\textit{true} \text{ orelse } \textit{false})) \text{ implies } \textit{false} \\ &= \textit{false} \end{aligned}$$



# Structural Induction

A technique for proving properties about inductively defined sets.

To prove that a proposition  $P(s)$  is true for all structures  $s$ , prove the following:

- 1 (Base case)  $P$  is true on simple structures (those without substructures)
- 2 (Inductive case) If  $P$  is true on the substructures of  $s$ , then it is true on  $s$  itself. The assumption is called *induction hypothesis (I.H.)*.

## Example 1

Let  $S$  be the set defined by the following inference rules:

$$\frac{}{\mathbf{3}} \quad \frac{x \quad y}{x + y}$$

Prove that for all  $x \in S$ ,  $x$  is divisible by  $\mathbf{3}$ .

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- (Inductive case) The induction hypothesis (I.H.) is

$$x \text{ is divisible by } \mathbf{3}, \quad y \text{ is divisible by } \mathbf{3}.$$

$$\text{Let } x = \mathbf{3}k_1 \text{ and } y = \mathbf{3}k_2.$$

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Let  $x = \mathbf{3}k_1$  and  $y = \mathbf{3}k_2$ . Using I.H., we derive

$$x + y \text{ is divisible by } \mathbf{3}$$

as follows:

$$\begin{aligned} x + y &= \mathbf{3}k_1 + \mathbf{3}k_2 \quad \dots \text{ by I.H.} \\ &= \mathbf{3}(k_1 + k_2) \end{aligned}$$



## Example 2

Let  $S$  be the set defined by the following inference rules:

$$\frac{}{()} \quad \frac{x}{(x)} \quad \frac{x \quad y}{xy}$$

Prove that every element of the set has the same number of ( 's and ) 's.

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**Proof** Restate the claim formally:

$$\text{If } x \in S \text{ then } l(x) = r(x)$$

where  $l(x)$  and  $r(x)$  denote the number of ( 's and ) 's, respectively.

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We prove it by structural induction:

- (Base case): The base case is when  $x = ()$ . Then  $l(x) = 1 = r(x)$ .



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- (Inductive case): There are two inductive cases:

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Induction hypotheses (I.H.):

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- ▶ The first case. We prove  $l((x)) = r((x))$ :

$$\begin{aligned} l((x)) &= l(x) + 1 && \dots \text{by definition of } l((x)) \\ &= r(x) + 1 && \dots \text{by I.H.} \\ &= r((x)) && \dots \text{by definition of } r((x)) \end{aligned}$$

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- ▶ The second case. We prove  $l(xy) = r(xy)$ :

$$\begin{aligned} l(xy) &= l(x) + l(y) && \dots \text{by definition of } l(xy) \\ &= r(x) + r(y) && \dots \text{by I.H.} \\ &= r(xy) && \dots \text{by definition of } r(xy) \end{aligned}$$



## Example 3

Let  $\mathcal{T}$  be the set of binary trees:

$$\overline{\text{leaf}} \quad \frac{t_1 \quad t_2}{(n, t_1, t_2)} \quad n \in \mathbb{Z}$$

Prove that for all such trees, the number of leaves is always one more than the number of internal nodes.

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- (Base case): The base case is when  $t = \mathbf{leaf}$ , where  $l(t) = 1$  and  $i(t) = 0$ .

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Using I.H., we prove  $l((n, t_1, t_2)) = i((n, t_1, t_2)) + 1$ :

$$\begin{aligned} l((n, t_1, t_2)) &= l(t_1) + l(t_2) \\ &= i(t_1) + 1 + i(t_2) + 1 \quad \text{by induction hypothesis} \\ &= i(n, t_1, t_2) + 1 \end{aligned}$$



# Summary

- Computer science is full of inductive definitions.
  - ▶ primitive values: booleans, characters, integers, strings, etc
  - ▶ compound values: lists, trees, graphs, etc
  - ▶ language syntax and semantics
- Structural induction
  - ▶ a general technique for reasoning about inductively defined sets