

COSE312: Compilers

Lecture 1 — Overview of Compilers

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What is Compiler?

Software that translates a program written in one language (“source language”) into a program in another language (“target language”):



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Typically,

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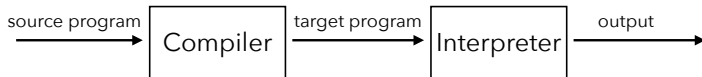
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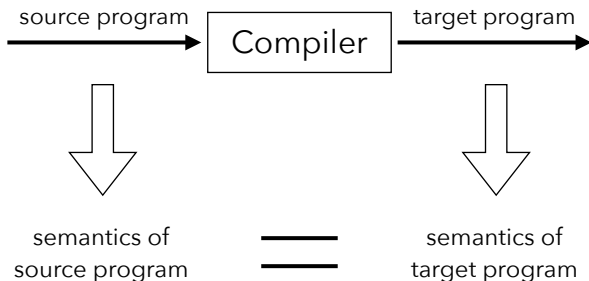
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The target language may not be a machine language:



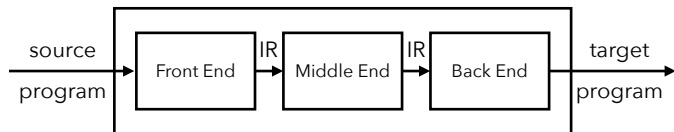
A Fundamental Requirement

- The compiler must preserve the meaning of the source program.



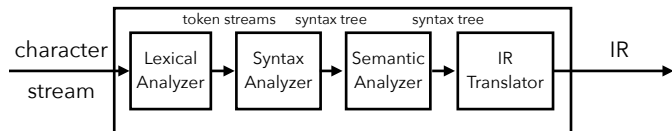
- Correctness of real-world compilers?

Structure of Compilers



- The front-end understands the source program and translates it to an intermediate representation (IR).
- The middle-end takes a program in IR and optimizes it in terms of efficiency, energy consumption, and so on.
- The back-end transforms the IR program into machine-code.

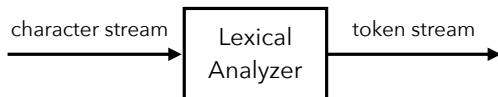
Front End



- Lexical analyzer transforms character streams into token streams.
- Syntax analyzer transforms token streams into syntax trees.
- Semantic analyzer checks the correctness of input programs.
- IR translator converts syntax trees into IRs.

Lexical¹ Analyzer (Lexer)

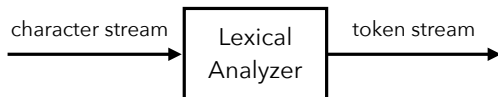
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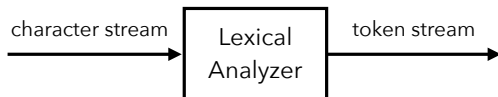
ex) The lexical analyzer for C transforms the character stream

```
pos = init + rate * 10
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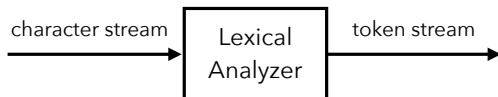
into the following sequence:

```
"pos", "=", "init", "+", "rate", "*", "10"
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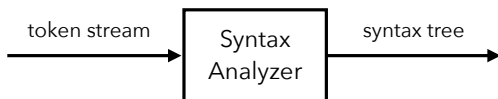
and then produces a *token* sequence:

```
(ID, pos), ASSIGN, (ID, init), PLUS, (ID, rate), MULT, (NUM,10)
```

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Syntax² Analyzer (Parser)

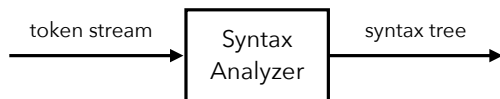
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²the way in which words are put together to form phrases, clauses, or sentences

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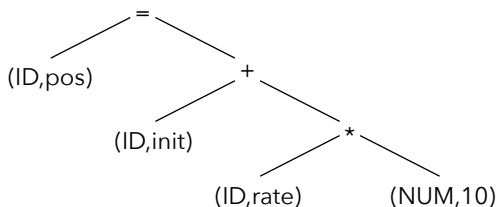
A parser analyzes the grammatical structure of the source program:



ex) A C parser transforms the sequence of tokens

(ID, pos), =, (ID, init), +, (ID, rate), *, (NUM,10)

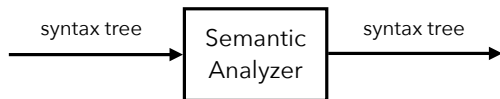
into the syntax tree:



²the way in which words are put together to form phrases, clauses, or sentences

Semantic Analyzer

A semantic analyzer aims to establish the correctness of the input program:



- Safety errors

- ▶ type errors, e.g.,

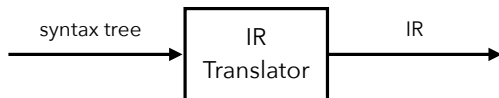
```
int x = 1;  
string y = "hello";  
int z = x + y;
```

- ▶ memory errors (array out of bounds, null-dereference, etc)

- Functional errors

- ▶ pre/post conditions

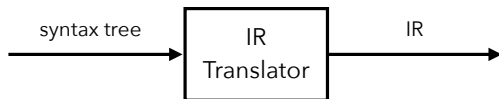
IR Translator



Intermediate Representation:

- lower-level than the source language
- higher-level than the target language

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ex) Three-address code:

```
t1 = 10
```

```
t2 = rate * t1
```

```
t3 = init + t2
```

```
pos = t3
```

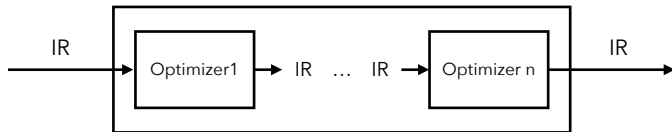

Optimizer

Transforms IR to have better performance:



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E.g.,

```
t1 = 10
t2 = rate * t1
t3 = init + t2
pos = t3
```

original IR

```
t1 = 10
t2 = rate * 10
t3 = init + t2
pos = t3
```

```
t2 = rate * 10
t3 = init + t2
pos = t3
```

```
t2 = rate * 10
pos = init + t2
```

final IR

Back End

Generates the target machine code:



ex) From the IR

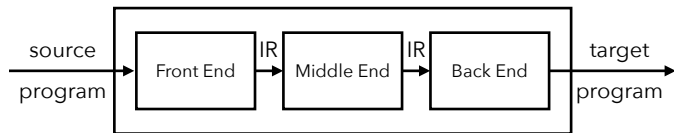
```
t2 = rate * 10  
pos = init + t2
```

generates the machine code

```
LOAD  R2, rate  
MUL   R2, R2, #10  
LOAD  R1, init  
ADD   R1, R1, R2  
STORE pos, R1
```

Summary

Compilers consist of three phases:



- Front end understands the syntax and semantics of source program.
- Middle end improves the efficiency of the program.
- Back end generates the target program.