COSE215: Theory of Computation

Lecture 4 — Regular Expressions

Hakjoo Oh 2016 Spring

Regular expression

A regular expression denotes a language.

E.g.,
$$(a+(b\cdot c))^*$$
 stands for:

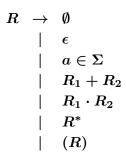
$$\{\epsilon, a, bc, aa, abc, bca, bcbc, aaa, aabc, \ldots\}$$

Syntax

Definition (Syntax of regular expressions)

Regular expressions over alphabet Σ are constructed recursively:

- **1** (Basis) \emptyset , ϵ , and $a \in \Sigma$ are regular expressions.
- (Induction)
 - If R_1 and R_2 are regular expressions, so are R_1+R_2 and $R_1\cdot R_2$.
 - If R is a regular expression, so are R^* and (R).



Semantics

Definition (Semantics of regular expressions)

A regular expression $m{R}$ means a set of strings, denoted $m{L}(m{R})$, which is defined inductively:

$$L(\emptyset) = \emptyset$$
 $L(\epsilon) = \{\epsilon\}$
 $L(a) = \{a\}$
 $L(R_1 + R_2) = L(R_1) \cup L(R_2)$
 $L(R_1 \cdot R_2) = L(R_1)L(R_2)$
 $L(R^*) = (L(R))^*$
 $L(R) = L(R)$

Example

$$L(a^* \cdot (a+b)) =$$

Exercises

Find the languages of the regular expressions and equivalent finite automata.

- $(a+b)^*$
- $(a+b)^*(a+b)$
- $\bullet \ (a \cdot a)^*(b \cdot b)^*b$

Exercises

Find regular expressions for the languages:

- ullet $L=\{w\in\{0,1\}^*\mid ext{0} ext{ and } ext{1} ext{ alternate in } w\}$
- ullet $L=\{w\in\{0,1\}^*\mid w ext{ has at least one pair of consecutive zeros}\}$
- $\bullet \ L = \{a^nb^m \mid n \geq 3, m \text{ is even}\}$
- $\bullet \ L = \{a^nb^m \mid (n+m) \text{ is even}\}$
- $\bullet \ L = \{a^nb^m \mid n \geq 4, m \leq 3\}$

Memo

Memo