# COSE212: Programming Languages

Lecture 13 — Automatic Type Inference (1)

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# The Problem of Automatic Type Inference

Given a program E, infer the most general type of E if E can be typed (i.e.,  $[] \vdash E : t$  for some  $t \in T$ ). If E cannot be typed, say so.

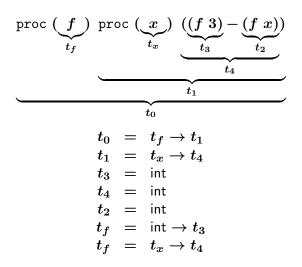
- let  $f = \operatorname{proc}(x)(x+1)$  in  $(\operatorname{proc}(x)(x1)) f$
- let f = proc (x) (x + 1) in (proc (x) (x true)) f
- ullet proc (x) x

# Automatic Type Inference

- A static analysis algorithm that automatically figures out types of expressions by observing how they are used.
- The algorithm is *sound and complete* with respect to the type system design.
  - (Sound) If the analysis finds a type for an expression, the expression is well-typed with the type according to the type system.
  - (Complete) If an expression has a type according to the type system, the analysis is guaranteed to find the type.
- The algorithm consists of two steps:
  - Generate type equations from the program text.
  - Solve the equations.

#### Generating Type Equations

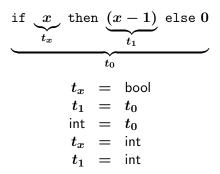
For every subexpression and variable, introduce type variables and derive equations between the type variables.

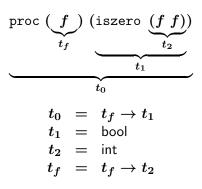


$$\underbrace{\frac{f}{t_f}\underbrace{(f\ 11)}_{t_0}}_{t_0}$$

$$\underbrace{t_0}_{t_0} = t_f \to t_1$$

$$t_f = \operatorname{int} \to t_1$$





# Idea: Deriving Equations from Typing Rules

For each expression e and variable x, let  $t_e$  and  $t_x$  denote the type of the expression and variable. Then, the typing rules dictate the equations that must hold between the type variables.

$$egin{aligned} rac{\Gamma dash E_1: \mathsf{int} & \Gamma dash E_2: \mathsf{int}}{\Gamma dash E_1 + E_2: \mathsf{int}} \ & t_{E_1} = \mathsf{int} \ \land \ t_{E_2} = \mathsf{int} \ \land \ t_{E_1 + E_2} = \mathsf{int} \end{aligned}$$

$$egin{array}{c} \Gamma dash E : \mathsf{int} \ \Gamma dash \mathsf{iszero} \ E : \mathsf{bool} \end{array}$$

$$t_E = \mathsf{int} \ \land \ t_{(\mathsf{iszero}\ E)} = \mathsf{bool}$$

$$egin{aligned} egin{aligned} rac{\Gamma dash E_1:t_1 
ightarrow t_2 & \Gamma dash E_2:t_1 \ & \Gamma dash E_1 E_2:t_2 \end{aligned}}{t_{E_1} = t_{E_2} 
ightarrow t_{(E_1 \ E_2)}} \end{aligned}$$

# Idea: Deriving Equations from Typing Rules

$$\begin{array}{lll} \Gamma \vdash E_1 : \mathsf{bool} & \Gamma \vdash E_2 : t & \Gamma \vdash E_3 : t \\ \hline \Gamma \vdash \mathsf{if} \ E_1 \ \mathsf{then} \ E_2 \ \mathsf{else} \ E_3 : t \\ \\ t_{E_1} &=& \mathsf{bool} \ \land \\ t_{E_2} &=& t_{(\mathsf{if} \ E_1 \ \mathsf{then} \ E_2 \ \mathsf{else} \ E_3)} \ \land \\ t_{E_3} &=& t_{(\mathsf{if} \ E_1 \ \mathsf{then} \ E_2 \ \mathsf{else} \ E_3)} \\ \bullet & \hline \frac{[x \mapsto t_1]\Gamma \vdash E : t_2}{\Gamma \vdash \mathsf{proc} \ x \ E : t_1 \to t_2} \\ \bullet & \hline t_{(\mathsf{proc} \ (x) \ E)} = t_x \to t_E \\ \bullet & \hline \Gamma \vdash E_1 : t_1 \quad [x \mapsto t_1]\Gamma \vdash E_2 : t_2 \\ \bullet & \hline \Gamma \vdash \mathsf{let} \ x = E_1 \ \mathsf{in} \ E_2 : t_2 \\ \hline t_x &=& t_{E_1} \ \land \ t_{E_2} = t_{(\mathsf{let} \ x = E_1 \ \mathsf{in} \ E_2)} \end{array}$$

#### Summary

The algorithm for automatic type inference:

- Generate type equations from the program text.
  - Introduce type variables for each subexpression and variable.
  - ► Generate equations between type variables according to typing rules.
- Solve the equations.