COSE212: Programming Languages

Lecture 6 — Design and Implementation of PLs (2) Procedures

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2017 Fall

Review: The Let Language

Syntax:

$$\begin{array}{rcl} P & \rightarrow & E \\ E & \rightarrow & n \\ & | & x \\ & | & E + E \\ & | & E - E \\ & | & \texttt{iszero } E \\ & | & \texttt{if } E \texttt{ then } E \texttt{ else } E \\ & | & \texttt{let } x = E \texttt{ in } E \\ & | & \texttt{read} \end{array}$$

Review: The Let Language

Semantic domain:

$$\begin{array}{rcl} \mathit{Val} & = & \mathbb{Z} + \mathit{Bool} \\ \mathit{Env} & = & \mathit{Var} \to \mathit{Val} \end{array}$$

Semantics rules:

$$\frac{}{\rho \vdash n \Rightarrow n} \qquad \frac{}{\rho \vdash x \Rightarrow \rho(x)}$$

$$\frac{\rho \vdash E_1 \Rightarrow n_1 \qquad \rho \vdash E_2 \Rightarrow n_2}{\rho \vdash E_1 + E_2 \Rightarrow n_1 + n_2} \qquad \frac{\rho \vdash E_1 \Rightarrow n_1 \qquad \rho \vdash E_2 \Rightarrow n_2}{\rho \vdash E_1 - E_2 \Rightarrow n_1 - n_2}$$

$$\frac{}{\rho \vdash \texttt{read} \Rightarrow n} \qquad \frac{\rho \vdash E \Rightarrow 0}{\rho \vdash \texttt{iszero}\ E \Rightarrow \mathit{true}} \qquad \frac{\rho \vdash E \Rightarrow n}{\rho \vdash \texttt{iszero}\ E \Rightarrow \mathit{false}}\ n \neq 0$$

$$\frac{\rho \vdash E_1 \Rightarrow \mathit{true} \qquad \rho \vdash E_2 \Rightarrow v}{\rho \vdash \texttt{if}\ E_1\ \texttt{then}\ E_2\ \texttt{else}\ E_3 \Rightarrow v} \qquad \frac{\rho \vdash E_1 \Rightarrow \mathit{false} \qquad \rho \vdash E_3 \Rightarrow v}{\rho \vdash \texttt{if}\ E_1\ \texttt{then}\ E_2\ \texttt{else}\ E_3 \Rightarrow v}$$

$$\frac{\rho \vdash E_1 \Rightarrow v_1 \qquad [x \mapsto v_1]\rho \vdash E_2 \Rightarrow v}{\rho \vdash \texttt{let}\ x = E_1\ \texttt{in}\ E_2 \Rightarrow v}$$

Proc = Let + Procedures

$$
\begin{array}{rcl}
P & \rightarrow & E \\
E & \rightarrow & n \\
& | & x \\
& | & E + E \\
& | & E - E \\
& | & \texttt{iszero}\ E \\
& | & \texttt{if}\ E\ \texttt{then}\ E\ \texttt{else}\ E \\
& | & \texttt{let}\ x = E\ \texttt{in}\ E \\
& | & \texttt{read} \\
& | & \texttt{proc}\ x\ E \\
& | & E\ E
\end{array}
$$

Example

- ```
  let f = proc (x) (x-11)
  in (f (f 77))
  ```
  ```

Example

- `let f = proc (x) (x-11)`
  `in (f (f 77))`
- `((proc (f) (f (f 77))) (proc (x) (x-11)))`

Free/Bound Variables of Procedures

- An occurrence of the variable `x` is *bound* when it occurs without definitions in the body of a procedure whose formal parameter is `x`.
- Otherwise, the variable is *free*.
- Examples:
  - `proc (y) (x+y)`
  - `proc (x) (let y = 1 in x + y + z)`
  - `proc (x) (proc (y) (x+y))`
  - `let x = 1 in proc (y) (x+y)`
  - `let x = 1 in proc (y) (x+y+z)`

Static vs. Dynamic Scoping

What is the result of the program?

```
let x = 1
in let f = proc (y) (x+y)
   in let x = 2
      in let g = proc (y) (x+y)
         in (f 1) + (g 1)
```

Static vs. Dynamic Scoping

What is the result of the program?

```
let x = 1
in let f = proc (y) (x+y)
   in let x = 2
      in let g = proc (y) (x+y)
         in (f 1) + (g 1)
```

Two ways to determine free variables of procedures:

- In *static scoping* (*lexical scoping*), the procedure body is evaluated in the environment where the procedure is defined (i.e. procedure-creation environment).
- In *dynamic scoping*, the procedure body is evaluated in the environment where the procedure is called (i.e. calling environment)

Exercises

What is the result of the program?

- In static scoping:
- In dynamic scoping:

1.
```
let a = 3
in let p = proc (z) a
   in let f = proc (x) (p 0)
      in let a = 5
         in (f 2)
```

2.
```
let a = 3
in let p = proc (z) a
   in let f = proc (a) (p 0)
      in let a = 5
         in (f 2)
```

Why Static Scoping?

Most modern languages use static scoping. Why?

- Reasoning about programs is much simpler in static scoping.
- In static scoping, renaming bound variables by their lexical definitions does not change the semantics, which is unsafe in dynamic scoping.

```
let x = 1
in let f = proc (y) (x+y)
   in let x = 2
      in let g = proc (y) (x+y)
         in (f 1) + (g 1)
```

- In static scoping, names are resolved at compile-time.
- In dynamic scoping, names are resolved only at runtime.

Semantics of Procedures: Static Scoping

- Domain:

$$\begin{array}{rcl} \boldsymbol{Val} & = & \mathbb{Z} + \boldsymbol{Bool} + \boldsymbol{Procedure} \\ \boldsymbol{Procedure} & = & \boldsymbol{Var} \times \boldsymbol{E} \times \boldsymbol{Env} \\ \boldsymbol{Env} & = & \boldsymbol{Var} \to \boldsymbol{Val} \end{array}$$

  The procedure value is called *closures*. The procedure is closed in its creation environment.

Semantics of Procedures: Static Scoping

- Domain:

$$\begin{array}{rcl} \mathit{Val} & = & \mathbb{Z} + \mathit{Bool} + \mathit{Procedure} \\ \mathit{Procedure} & = & \mathit{Var} \times E \times \mathit{Env} \\ \mathit{Env} & = & \mathit{Var} \to \mathit{Val} \end{array}$$

  The procedure value is called *closures*. The procedure is closed in its creation environment.

- Semantics rules:

$$\frac{}{\rho \vdash \texttt{proc}\ x\ E \Rightarrow (x, E, \rho)}$$

$$\frac{\rho \vdash E_1 \Rightarrow (x, E, \rho') \qquad \rho \vdash E_2 \Rightarrow v \qquad [x \mapsto v]\rho' \vdash E \Rightarrow v'}{\rho \vdash E_1\ E_2 \Rightarrow v'}$$

Examples

$$\frac{}{[] \vdash \texttt{(proc (x) (x)) 1} \Rightarrow \mathbf{1}}$$

Examples

$$\dfrac{}{[] \vdash \begin{array}{l}\texttt{let x = 1} \\ \texttt{in let f = proc (y) (x+y)} \\ \quad\texttt{in let x = 2} \\ \qquad\texttt{in (f 3)}\end{array} \Rightarrow \mathbf{4}}$$

Semantics of Procedures: Dynamic Scoping

- Domain:

$$\begin{array}{rcl} \boldsymbol{Val} & = & \mathbb{Z} + \boldsymbol{Bool} + \boldsymbol{Procedure} \\ \boldsymbol{Procedure} & = & \boldsymbol{Var} \times \boldsymbol{E} \\ \boldsymbol{Env} & = & \boldsymbol{Var} \rightarrow \boldsymbol{Val} \end{array}$$

- Semantics rules:

$$\frac{}{\rho \vdash \texttt{proc}\ x\ E \Rightarrow (x, E)}$$

$$\frac{\rho \vdash E_1 \vdash (x, E) \qquad \rho \vdash E_2 \Rightarrow v \qquad [x \mapsto v]\rho \vdash E \Rightarrow v'}{\rho \vdash E_1\ E_2 \Rightarrow v'}$$

Examples

$$\frac{}{[\,] \vdash \begin{array}{l}\texttt{let x = 1}\\ \texttt{in let f = proc (y) (x+y)}\\ \quad\texttt{in let x = 2}\\ \qquad\texttt{in (f 3)}\end{array} \Rightarrow \mathbf{5}}$$

cf) Multiple Argument Procedures

- We can get the effect of multiple argument procedures by using procedures that return other procedures.
- ex) a function that takes two arguments and return their sum:

```
let f = proc (x) proc (y) (x+y)
in ((f 3) 4)
```

Adding Recursive Procedures

The current language does not support recursive procedures, e.g.,

```
let f = proc (x) (f x)
in (f 1)
```

for which evaluation gets stuck:

$$\dfrac{[f \mapsto (x, f\ x, [])] \vdash f \Rightarrow (x, f\ x, []) \qquad \dfrac{[x \mapsto 1] \vdash f \Rightarrow ? \qquad [x \mapsto 1] \vdash x \Rightarrow 1}{[x \mapsto 1] \vdash f\ x \Rightarrow ?}}{[f \mapsto (x, f\ x, [])] \vdash (f\ 1) \Rightarrow ?}$$

Two solutions:

- go back to dynamic scoping :-(
- modify the language syntax and semantics for procedure :-)

Recursion is Not Special in Dynamic Scoping

With dynamic scoping, recursive procedures require no special mechanism. Running the program

```
let f = proc (x) (f x)
in (f 1)
```

via dynamic scoping semantics

$$\frac{\rho \vdash E_1 \vdash (x, E) \qquad \rho \vdash E_2 \Rightarrow v \qquad [x \mapsto v]\rho \vdash E \Rightarrow v'}{\rho \vdash E_1\ E_2 \Rightarrow v'}$$

proceeds well:

$$\cfrac{\cfrac{\cfrac{\cfrac{\vdots}{[f \mapsto (x, f\ x), x \mapsto 1] \vdash \texttt{f x} \Rightarrow}}{[f \mapsto (x, f\ x), x \mapsto 1] \vdash \texttt{f x} \Rightarrow}}{[f \mapsto (x, f\ x)] \vdash \texttt{f 1} \Rightarrow}}{[] \vdash \texttt{let f = proc (x) (f x) in (f 1)} \Rightarrow}$$

Adding Recursive Procedures

$$
\begin{array}{rcl}
P & \rightarrow & E \\
E & \rightarrow & n \\
& | & x \\
& | & E + E \\
& | & E - E \\
& | & \texttt{iszero } E \\
& | & \texttt{if } E \texttt{ then } E \texttt{ else } E \\
& | & \texttt{let } x = E \texttt{ in } E \\
& | & \texttt{read} \\
& | & \texttt{letrec } f(x) = E \texttt{ in } E \\
& | & \texttt{proc } x \; E \\
& | & E \; E
\end{array}
$$

Example

```
letrec double(x) =
  if zero?(x) then 0 else ((double (x-1)) + 2)
in (double 1)
```

Semantics of Recursive Procedures

- Domain:

$$\begin{array}{rcl} \boldsymbol{Val} & = & \mathbb{Z} + \boldsymbol{Bool} + \boldsymbol{Procedure} + \boldsymbol{RecProcedure} \\ \boldsymbol{Procedure} & = & \boldsymbol{Var} \times \boldsymbol{E} \times \boldsymbol{Env} \\ \boldsymbol{RecProcedure} & = & \boldsymbol{Var} \times \boldsymbol{Var} \times \boldsymbol{E} \times \boldsymbol{Env} \\ \boldsymbol{Env} & = & \boldsymbol{Var} \to \boldsymbol{Val} \end{array}$$

- Semantics rules:

$$\frac{[f \mapsto (f, x, E_1, \rho)]\rho \vdash E_2 \Rightarrow v}{\rho \vdash \texttt{letrec } f(x) = E_1 \texttt{ in } E_2 \Rightarrow v}$$

$$\frac{\begin{array}{c} \rho \vdash E_1 \Rightarrow (f, x, E, \rho') \quad \rho \vdash E_2 \Rightarrow v \\ [x \mapsto v, f \mapsto (f, x, E, \rho')]\rho' \vdash E \Rightarrow v' \end{array}}{\rho \vdash E_1\ E_2 \Rightarrow v'}$$

Example

$$
\dfrac{[f \mapsto (f, x, f\ x, [])] \vdash f \Rightarrow (f, x, f\ x, []) \qquad \dfrac{\dfrac{\vdots}{[x \mapsto 1, f \mapsto (f, x, f\ x, [])] \vdash f\ x \Rightarrow}}{[x \mapsto 1, f \mapsto (f, x, f\ x, [])] \vdash f\ x \Rightarrow}}{\dfrac{[f \mapsto (f, x, f\ x, [])] \vdash f\ 1 \Rightarrow}{[] \vdash \texttt{letrec}\ f(x) = f\ x\ \texttt{in}\ f\ 1 \Rightarrow}}
$$

Summary: The Proc Language

A programming language with expressions and procedures:

Syntax

$$
\begin{array}{rcl}
P & \to & E \\
E & \to & n \\
& \mid & x \\
& \mid & E + E \\
& \mid & E - E \\
& \mid & \texttt{iszero}\ E \\
& \mid & \texttt{if}\ E\ \texttt{then}\ E\ \texttt{else}\ E \\
& \mid & \texttt{let}\ x = E\ \texttt{in}\ E \\
& \mid & \texttt{read} \\
& \mid & \texttt{letrec}\ f(x) = E\ \texttt{in}\ E \\
& \mid & \texttt{proc}\ x\ E \\
& \mid & E\ E
\end{array}
$$

Summary

Semantics

$$\frac{}{\rho \vdash n \Rightarrow n} \qquad \frac{}{\rho \vdash x \Rightarrow \rho(x)} \qquad \frac{\rho \vdash E_1 \Rightarrow n_1 \qquad \rho \vdash E_2 \Rightarrow n_2}{\rho \vdash E_1 + E_2 \Rightarrow n_1 + n_2}$$

$$\frac{\rho \vdash E \Rightarrow 0}{\rho \vdash \texttt{iszero}\ E \Rightarrow \mathit{true}} \qquad \frac{\rho \vdash E \Rightarrow n}{\rho \vdash \texttt{iszero}\ E \Rightarrow \mathit{false}}\ n \neq 0 \qquad \frac{}{\rho \vdash \texttt{read} \Rightarrow n}$$

$$\frac{\rho \vdash E_1 \Rightarrow \mathit{true} \qquad \rho \vdash E_2 \Rightarrow v}{\rho \vdash \texttt{if}\ E_1\ \texttt{then}\ E_2\ \texttt{else}\ E_3 \Rightarrow v} \qquad \frac{\rho \vdash E_1 \Rightarrow \mathit{false} \qquad \rho \vdash E_3 \Rightarrow v}{\rho \vdash \texttt{if}\ E_1\ \texttt{then}\ E_2\ \texttt{else}\ E_3 \Rightarrow v}$$

$$\frac{\rho \vdash E_1 \Rightarrow v_1 \qquad [x \mapsto v_1]\rho \vdash E_2 \Rightarrow v}{\rho \vdash \texttt{let}\ x = E_1\ \texttt{in}\ E_2 \Rightarrow v} \qquad \frac{[f \mapsto (f, x, E_1, \rho)]\rho \vdash E_2 \Rightarrow v}{\rho \vdash \texttt{letrec}\ f(x) = E_1\ \texttt{in}\ E_2 \Rightarrow v}$$

$$\frac{}{\rho \vdash \texttt{proc}\ x\ E \Rightarrow (x, E, \rho)}$$

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$$\frac{\rho \vdash E_1 \Rightarrow (f, x, E, \rho') \quad \rho \vdash E_2 \Rightarrow v \quad [x \mapsto v, f \mapsto (f, x, E, \rho')]\rho' \vdash E \Rightarrow v'}{\rho \vdash E_1\ E_2 \Rightarrow v'}$$

Mid-term

- Homework 2 will replace mid-term exam.
- No class on 10/24(Tue) and 10/26(Thr).