

COSE212: Programming Languages

Lecture 1 — Inductive Definitions (1)

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Inductive Definitions

A technique for formally defining a set:

- The set is defined in terms of itself.
- The only way of defining an infinite set by a finite means.

Three styles to define a set inductively:

- Top-down
- Bottom-up
- Rules of inference

Example (Top-Down)

Definition (S)

A natural number n is in S if and only if

- 1 $n = 0$, or
- 2 $n - 3 \in S$.

- Inductive definition of a set of natural numbers: $S \subseteq \mathbb{N} = \{0, 1, \dots\}$
- $\{0, 3, 6, 9, \dots\} \subseteq S$
- $\{0, 3, 6, 9, \dots\} \supseteq S$

$$S = \{0, 3, 6, 9, \dots\}.$$

Formal Proofs

Lemma

$$\{0, 3, 6, 9, \dots\} \subseteq S$$

By induction. To show: $3k \in S$ for all $k \in \mathbb{N}$.

- 1 Base case: $3k \in S$ when $k = 0$.
- 2 Inductive case: Assume $3k \in S$ (Induction Hypothesis, I.H.).
To show is $3 \cdot (k + 1) \in S$, which holds because
 $3 \cdot (k + 1) - 3 = 3k \in S$ by the induction hypothesis.

Lemma

$$\{0, 3, 6, 9, \dots\} \supseteq S$$

By proof by contradiction. Let $n = 3k + q$ ($q = 1$ or 2) and assume $n \in S$. By the definition of S , $n - 3, n - 6, \dots, n - 3k \in S$. Thus, S must include 1 or 2 , a contradiction.

A Bottom-up Definition

Definition (\mathcal{S})

\mathcal{S} is the *smallest* set such that $\mathcal{S} \subseteq \mathbb{N}$ and \mathcal{S} satisfies the following two conditions:

- 1 $0 \in \mathcal{S}$, and
- 2 if $n \in \mathcal{S}$, then $n + 3 \in \mathcal{S}$.

- The two conditions imply $\{0, 3, 6, 9, \dots\} \subseteq \mathcal{S}$.
- The two conditions do not imply $\{0, 3, 6, 9, \dots\} \supseteq \mathcal{S}$, e.g., \mathbb{N} .
- By requiring \mathcal{S} to be the **smallest** such a set,

$$\mathcal{S} = \{0, 3, 6, 9, \dots\}.$$

- The smallest set satisfying the conditions is unique.
 - ▶ Proof) If \mathcal{S}_1 and \mathcal{S}_2 satisfy the conditions and are both smallest, then $\mathcal{S}_1 \subseteq \mathcal{S}_2$ and $\mathcal{S}_2 \subseteq \mathcal{S}_1$. Therefore, $\mathcal{S}_1 = \mathcal{S}_2$ (\subseteq is anti-symmetric).

Rules of Inference

$$\frac{A}{B}$$

- A : hypothesis (antecedent)
- B : conclusion (consequent)
- “if A is true then B is also true”.
- \overline{B} : axiom.

Defining a Set by Rules of Inferences

Definition

$$\overline{0 \in S}$$
$$\frac{n \in S}{(n + 3) \in S}$$

Interpret the rules as follows:

“A natural number n is in S iff $n \in S$ can be derived from the axiom by applying the inference rules finitely many times”

ex) $3 \in S$ because

$$\overline{0 \in S} \text{ the axiom}$$
$$\frac{0 \in S}{3 \in S} \text{ the second rule}$$

Note that this interpretation enforces that S is the smallest set closed under the inference rules.

Summary

In inductive definitions, a set is defined in terms of itself. Three styles:

- Top-down
- Bottom-up
- Rules of inference

In PL, we mainly use the rules-of-inference method.

Exercises

- ① What set is defined by the following inductive rules?

$$\overline{\mathbf{3}} \quad \frac{x \quad y}{x + y}$$

- ② What set is defined by the following inductive rules?

$$\overline{()} \quad \frac{s}{(s)} \quad \frac{s}{s \ s}$$

- ③ Define the following set as rules of inference:

$$S = \{a, b, aa, ab, ba, bb, aaa, aab, aba, abb, baa, bab, bba, bbb, \dots\}$$

- ④ Define the following set as rules of inference:

$$S = \{x^n y^{n+1} \mid n \in \mathbb{N}\}$$