COSE215: Theory of Computation

Lecture 12 — Pushdown Automata (2)

Hakjoo Oh
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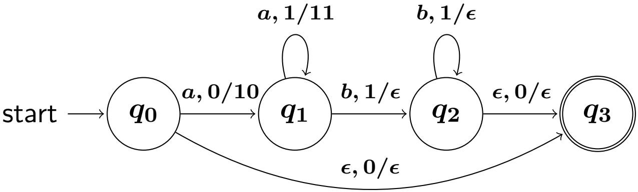
Exercises

- $L = \{a^n b^n \mid n \geq 0\}$:

$$P = (\{q_0, q_1, q_2, q_3\}, \{a, b\}, \{0, 1\}, \delta, q_0, 0, \{q_3\})$$



- $L = \{w \in \{a,b\}^* \mid n_a(w) = n_b(w)\}$
  - Intuition: Whenever we read a, we insert a counter symbol 0 onto the stack, and pop one counter symbol from the stack whenever b is found. For example,

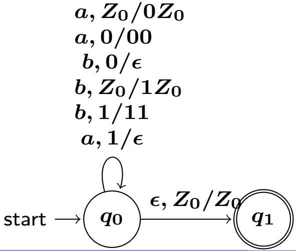
$$(abab, Z_0) \rightarrow (bab, 0Z_0) \rightarrow (ab, Z_0) \rightarrow (b, 0Z_0) \rightarrow (\epsilon, Z_0)$$

  - For the cases where a prefix of the input string contains more b's than a's, use a negative counter symbol, say 1, for counting the b's that should be matched against a's later. For example,

$$(bbaa, Z_0) \rightarrow (baa, 1Z_0) \rightarrow (aa, 11Z_0) \rightarrow (a, 1Z_0) \rightarrow (\epsilon, Z_0)$$

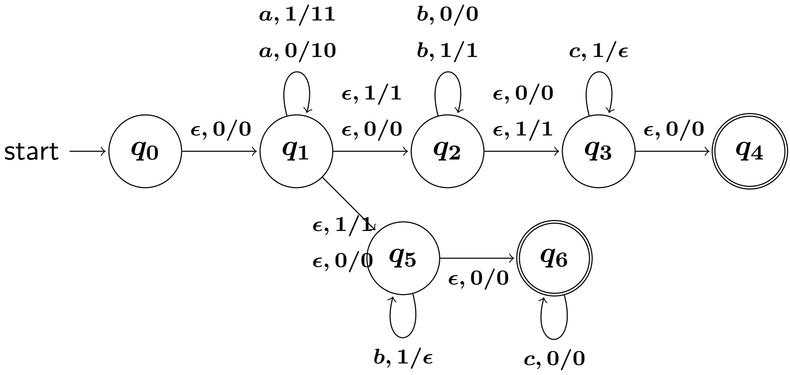
  - The pushdown automaton:

$$P = (\{q_0, q_1\}, \{a, b\}, \{0, 1, Z_0\}, \delta, q_0, Z_0, \{q_1\})$$



- $L = \{a^i b^j c^k \mid i,j,k \geq 0 \wedge (i = j \vee i = k)\}$
  - Think of the two cases separately:
    1. $L_1 = \{a^i b^j c^k \mid i,j,k \geq 0 \wedge i = j\}$.
    2. $L_2 = \{a^i b^j c^k \mid i,j,k \geq 0 \wedge j = k\}$.

$$P = (\{q_0, q_1, q_2, q_3, q_4, q_5, q_6\}, \{a, b, c\}, \{0, 1\}, \delta, q_0, 0, \{q_4, q_6\})$$



Configurations of PDA

- A configuration of a PDA consists of the automaton state and the stack contents.
- The configuration or instantaneous description (ID) is represented by (q, w, γ), where
  - q is the state,
  - w is the remaining input, and
  - γ is the stack contents.
- Suppose $(q, aw, X\beta)$ is a configuration and $(p, \alpha) \in \delta(q, a, X)$. Then, the configuration moves in one step to $(p, w, \alpha\beta)$:

$$(q, aw, X\beta) \vdash (p, w, \alpha\beta)$$

The Language of Pushdown Automata

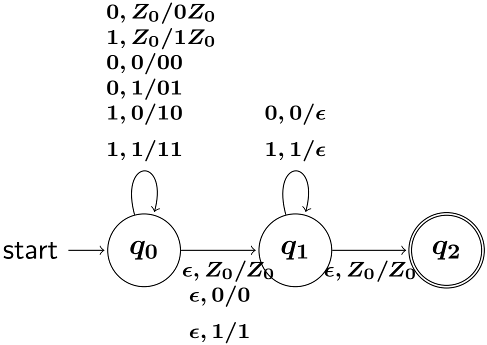
Definition (Acceptance by Final State)

Let $P = (Q, \Sigma, \Gamma, \delta, q_0, Z_0, F)$ be a PDA. Then $L(P)$, the language of P by final state, is

$$L(P) = \{w \in \Sigma^* \mid (q_0, w, Z_0) \vdash^* (q, \epsilon, \alpha)\}$$

for some state $q \in F$ and any stack string α.

Example



The PDA contains 1111, because $(q_0, 1111, Z_0) \vdash^* (q_2, \epsilon, Z_0)$.

Another Way of Defining The Language of a PDA

Definition (Acceptance by Empty Stack)

Let $P = (Q, \Sigma, \Gamma, \delta, q_0, Z_0, F)$ be a PDA. Then $N(P)$, the language of P accepted by empty stack, is

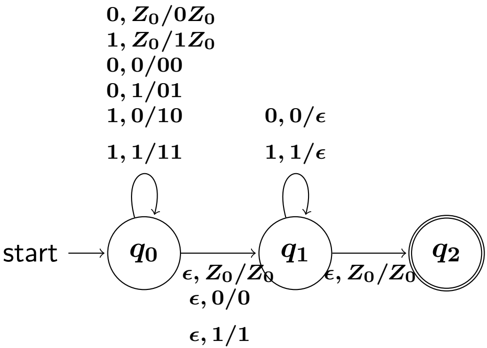
$$N(P) = \{w \in \Sigma^* \mid (q_0, w, Z_0) \vdash^* (q, \epsilon, \epsilon)\}$$

for any q.

When accepting by empty stack, we omit the F component:

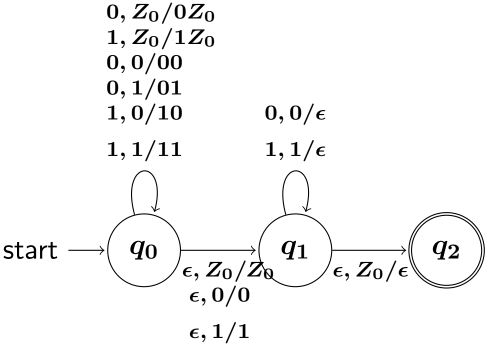
$$(Q, \Sigma, \Gamma, \delta, q_0, Z_0)$$

Example



- $L(P) =$
- $N(P) =$

Examples



- $L(P) =$
- $N(P) =$

Equivalence

Theorem (Equivalence of Final State and Empty Stack)

For any language L, there exists a PDA P_F such that $L = L(P_F)$ iff there exists a PDA P_N such that $L = N(P_N)$.

Lemma (From Empty Stack to Final State)

For any PDA $P_N = (Q, \Sigma, \Gamma, \delta_N, q_0, Z_0)$, there is a PDA P_F such that $N(P_N) = L(P_F)$.

Lemma (From Final State to Empty Stack)

For any PDA $P_F = (Q, \Sigma, \Gamma, \delta_F, q_0, Z_0, F)$, there is a PDA P_N such that $N(P_N) = L(P_F)$.

From Empty Stack to Final State

Given $P_N = (Q, \Sigma, \Gamma, \delta_N, q_0, Z_0)$, define

$$P_F = (Q \cup \{p_0, p_f\}, \Sigma, \Gamma \cup \{X_0\}, \delta_F, p_0, X_0, \{p_f\})$$

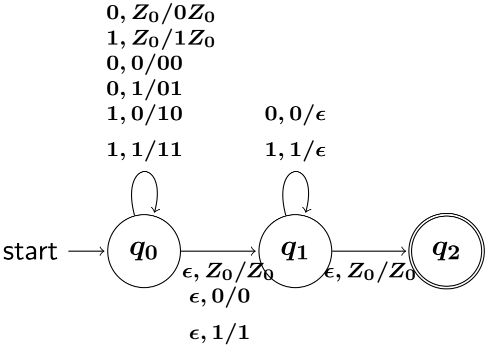
where

1. $\delta_F(p_0, \epsilon, X_0) = \{(q_0, Z_0X_0)\}$
2. For all $q \in Q$, $a \in \Sigma \cup \{\epsilon\}$, and $Y \in \Gamma$, $\delta_F(q, a, Y)$ contains $\delta_N(q, a, Y)$.
3. For all $q \in Q$, $\delta_F(q, \epsilon, X_0)$ contains (p_f, ϵ).

Then, w is in $L(P_F)$ if and only if w is in $N(P_N)$.

Example

Convert the following PDA to a PDA that accepts that same language by empty stack:



From Final State to Empty Stack

Given $P_F = (Q, \Sigma, \Gamma, \delta_F, q_0, Z_0, F)$, define

$$P_N = (Q \cup \{p_0, p\}, \Sigma, \Gamma \cup \{X_0\}, \delta_N, p_0, X_0)$$

where

1. $\delta_N(p_0, \epsilon, X_0) = \{(q_0, Z_0X_0)\}$
2. For all $q \in Q$, $a \in \Sigma \cup \{\epsilon\}$, and $Y \in \Gamma$, $\delta_N(q, a, Y)$ includes $\delta_F(q, a, Y)$.
3. For all accepting states $q \in F$ and $Y \in \Gamma \cup \{X_0\}$, $\delta_N(q, \epsilon, Y)$ includes (p, ϵ).
4. For all stack symbols $Y \in \Gamma \cup \{X_0\}$, $\delta_N(p, \epsilon, Y) = \{(p, \epsilon)\}$.

Equivalence of PDA's and CFG's

The following three classes of languages:

1. The context-free languages, i.e., the languages defined by CFG's.
2. The languages that are accepted by final state by some PDA.
3. The languages that are accepted by empty stack by some PDA.

are all the same class.

From CFG to PDA

Given a CFG $G = (V, T, P, S)$, define a PDA P (by empty stack):

$$P = (\{q\}, T, V \cup T, \delta, q, S)$$

where

- For each variable $A \in V$,

$$\delta(q, \epsilon, A) = \{(q, \beta) \mid (A \to \beta) \text{ is in } G\}$$

- For each terminal $a \in T$,

$$\delta(q, a, a) = \{(q, \epsilon)\}$$

Example

$$G = (\{B\}, \{(,)\}, P, B)$$

$$B \rightarrow BB \mid (B) \mid \epsilon$$

Deterministic Pushdown Automata

Definition

A pushdown automata $P = (Q, \Sigma, \Gamma, \delta, q_0, Z_0, F)$ is a *deterministic pushdown automata* (DPDA) if P makes at most one move at a time, i.e.,

1. $|\delta(q, a, X)| \leq 1$ for any $q \in Q$, $a \in \Sigma \cup \{\epsilon\}$, and $X \in \Gamma$.
2. If $\delta(q, a, X) \neq \emptyset$ for some $a \in \Sigma$, then $\delta(q, \epsilon, X) = \emptyset$.

Definition

A language L is said to be a deterministic context-free language iff there exists a DPDA P such that $L = L(P)$.

Example

The language

$$L = \{a^n b^n \mid n \geq 0\}$$

is a deterministic context-free language.

Fact1: DCFLs includes some CFLs

Example

The language

$$L = \{ww^R \mid w \in \{a,b\}^*\}$$

is *not* a deterministic context-free language.

Fact2: DCFLs do not include some CFLs

Regular Languages and DCFLs

Fact3: DCFLs include all RLs

Theorem

If L is a regular language, then $L = L(P)$ for some DPDA P.

Proof.

Let $A = (Q, \Sigma, \delta_A, q_0, F)$ be a DFA. Construct DPDA

$$P = (Q, \Sigma, \{Z_0\}, \delta_p, q_0, Z_0, F)$$

where define $\delta_p(q, a, Z_0) = \{(p, Z_0)\}$ for all p and q such that $\delta_A(q, a) = p$. Then, $(q_0, w, Z_0) \vdash^* (p, \epsilon, Z_0)$ iff $\delta_A^*(q_0, w) = p$. □

DPDA's and Ambiguous Grammars

Fact4: All DCFLs have unambiguous grammars.

Theorem

If $L = L(P)$ for some DPDA P, then L has an unambiguous grammar.

Fact5: DCFLs do not include all unambiguous CFLs.

The language

$$L = \{ww^R \mid w \in \{a, b\}^*\}$$

has an unambiguous grammar

$$S \to aSa \mid bSb \mid \epsilon$$

but not a DPDA language.

Summary

- PDA $=$ FA with a stack
- PDA is more powerful than FA. Cover *all* CFLs.
  - Still limited, e.g., $\{ww \mid w \in \Sigma^*\}$.
- DPDA is between FA and PDA

In general,

- FA with an external storage
  - queue, two stacks, random access memory, ...?
  - increase the language-recognizing power?