

Homework 5

COSE212, Fall 2019

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Due: 12/10, 23:59

Problem 1 Consider the language:

```
type exp =
  | CONST of int
  | VAR of var
  | ADD of exp * exp
  | SUB of exp * exp
  | MUL of exp * exp
  | DIV of exp * exp
  | READ
  | ISZERO of exp
  | IF of exp * exp * exp
  | LET of var * exp * exp
  | LETREC of var * var * exp * exp
  | PROC of var * exp
  | CALL of exp * exp
and var = string
```

Types for the language are defined as follows:

```
type typ = TyInt | TyBool | TyFun of typ * typ | TyVar of tyvar
and tyvar = string
```

Implement the following type-inference function:

`typeof : exp -> typ`

which takes a program and returns its type if the program is well-typed. When the program is ill-typed, `typeof` should raise an exception `TypeError`.

Examples:

- The program

```
PROC ("f",
  PROC ("x", SUB (CALL (VAR "f", CONST 3),
    CALL (VAR "f", VAR "x"))))
```

has type `TyFun (TyFun (TyInt, TyInt), TyFun (TyInt, TyInt))`.

- The program

```
PROC ("f", CALL (VAR "f", CONST 11))
```

has type `TyFun (TyFun (TyInt, TyVar "t"), TyVar "t")`, where `t` can be any type variable.

- The program

```
LET ("x", CONST 1,  
    IF (VAR "x", SUB (VAR "x", CONST 1), CONST 0))
```

is ill-typed, so `typeof` should raise an exception `TypeError`.

As discussed in class, `typeof` is defined with two functions: one for generating type equations and the other for solving the equations. Complete the implementation of these two functions:

```
gen_equations : TEnv.t -> exp -> typ -> typ_eqn  
solve         : typ_eqn -> Subst.t
```

Modules for type environments (`TEnv`) and substitutions (`Subst`), as well as the operations of applying substitutions to types (`Subst.apply`) and extending substitutions (`Subst.extend`), are provided.

Problem 2 Consider the language:

```
type exp =  
  | CONST of int  
  | VAR of var  
  | ADD of exp * exp  
  | SUB of exp * exp  
  | MUL of exp * exp  
  | DIV of exp * exp  
  | READ  
  | ISZERO of exp  
  | IF of exp * exp * exp  
  | LET of var * exp * exp  
  | LETREC of var * var * exp * exp  
  | PROC of var * exp  
  | CALL of exp * exp  
and var = string
```

Define the function

```
expand : exp -> exp
```

that transforms an expression into a semantically-equivalent expression where every let-bound variable in the original expression gets replaced by its definition. Examples and caveat:

- Evaluating

```
expand (LET ("x", CONST 1, VAR "x"))
```

produces CONST 1.

- Evaluating

```
expand (
  LET ("f", PROC ("x", VAR "x"),
    IF (CALL (VAR "f", ISZERO (CONST 0)),
      CALL (VAR "f", CONST 11),
      CALL (VAR "f", CONST 22))))
```

produces

```
IF (CALL (PROC ("x", VAR "x"), ISZERO (CONST 0)),
  CALL (PROC ("x", VAR "x"), CONST 11),
  CALL (PROC ("x", VAR "x"), CONST 22))
```

- Unused definitions should not go away. For example, evaluating

```
expand (LET ("x", ADD (CONST 1, ISZERO (CONST 0)), CONST 2))
```

should return LET ("x", ADD (CONST 1, ISZERO (CONST 0)), CONST 2),
not CONST 2.

As discussed in class, the function `expand` can be used for implementing the let-polymorphic type system. The type checker `typeof : exp -> typ` in Problem 1 does not support polymorphism and would not accept the program:

```
# typeof(
  LET ("f", PROC ("x", VAR "x"),
    IF (CALL (VAR "f", ISZERO (CONST 0)),
      CALL (VAR "f", CONST 11),
      CALL (VAR "f", CONST 22))));;
```

```
= Equations =
t2 = (t6 -> t7)
t7 = t6
(t5 -> bool) = t2
t5 = bool
int = int
(t4 -> t1) = t2
t4 = int
(t3 -> t1) = t2
t3 = int
```

The program does not have type. Rejected.

With `expand`, however, the same type checking algorithm will succeed:

```
# typeof(  
  expand(  
    LET ("f", PROC ("x", VAR "x"),  
      IF (CALL (VAR "f", ISZERO (CONST 0)),  
        CALL (VAR "f", CONST 11),  
        CALL (VAR "f", CONST 22)))));;
```

```
= Equations =  
(t8 -> bool) = (t9 -> t10)  
t10 = t9  
t8 = bool  
int = int  
(t5 -> t1) = (t6 -> t7)  
t7 = t6  
t5 = int  
(t2 -> t1) = (t3 -> t4)  
t4 = t3  
t2 = int
```

```
= Substitution =  
t3 |-> int  
t4 |-> int  
t2 |-> int  
t6 |-> int  
t7 |-> int  
t1 |-> int  
t5 |-> int  
t9 |-> bool  
t10 |-> bool  
t8 |-> bool
```

Type of the given program: int